



from tensor factorizations to circuits (and back)

lorenzo loconte

 @loreloc

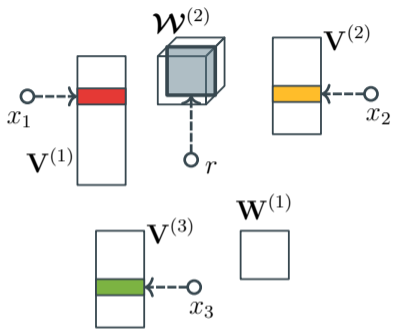
 @loreloc_

antonio vergari

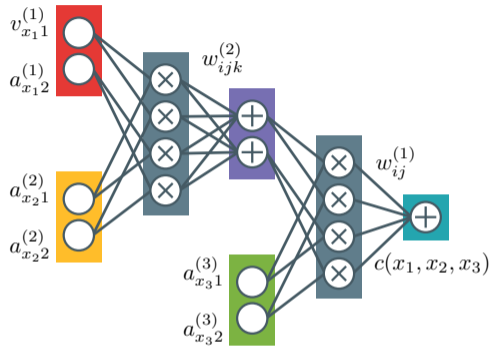
 @nolovedeep learning

 @tetraduazione

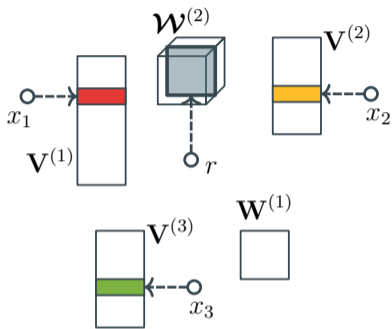
25th Feb 2025 – **AAAI-25** – Philadelphia



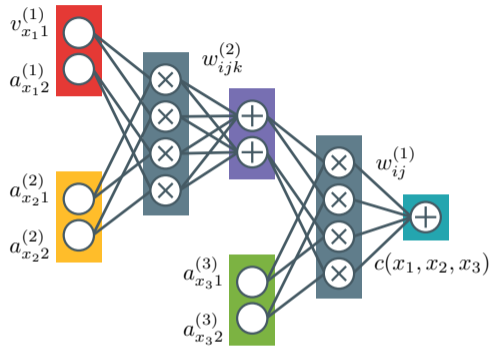
tensor factorizations



circuits

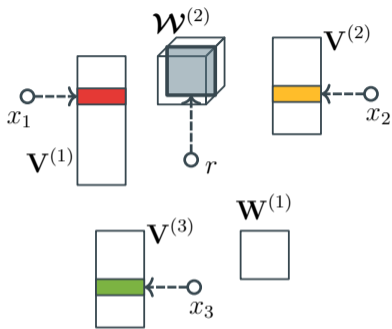


tensor factorizations

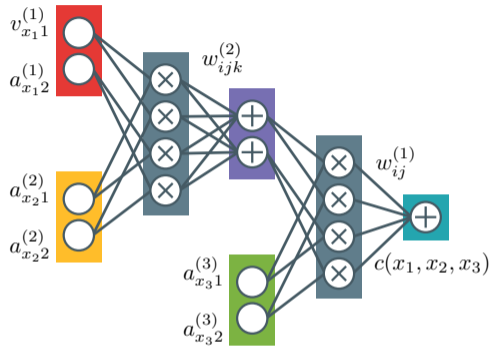


circuits

they look quite different...!?



tensor factorizations



circuits

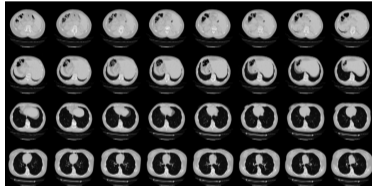
what can one take from another...!?

why tensor factorizations? (1/4)

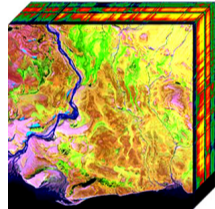
high-dimensional data as tensors



(PBS Nature)



(H. Zunair)



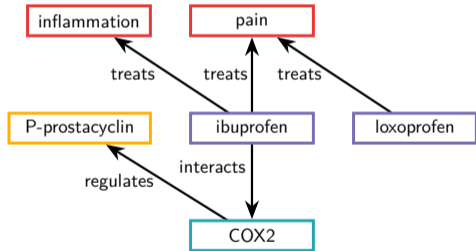
(N. M. Short)

Panagakis et al., "Tensor Methods in Computer Vision and Deep Learning", 2021

Wang et al., "Tensor Decompositions for Hyperspectral Data Processing in Remote Sensing: A Comprehensive Review", 2022

why tensor factorizations? (2/4)

graphs as tensors

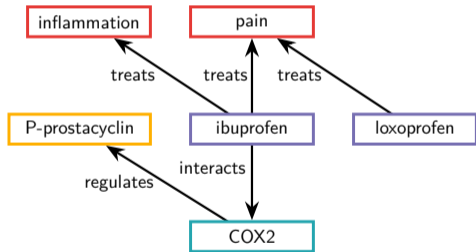


Entities:

- Drugs
- Symptoms
- Proteins
- Functions

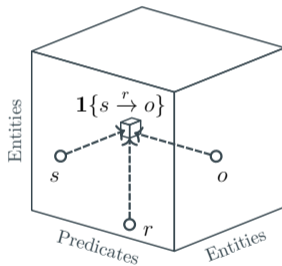
why tensor factorizations? (2/4)

graphs as tensors



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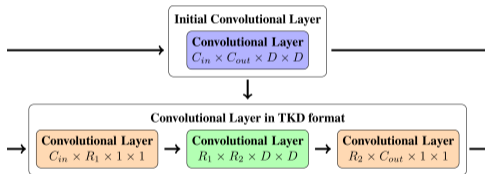


Knowledge base as Boolean tensor

[Nickel et al. 2016]

why tensor factorizations? (3/4)

compress ML models

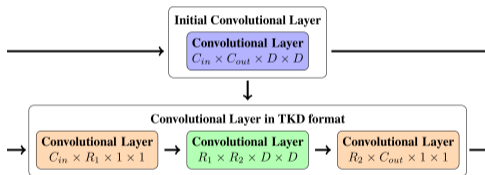


Compress convolutional layers

[Phan et al. 2020]

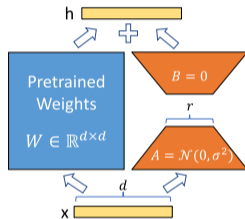
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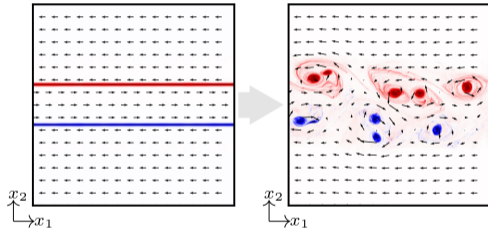


Low-rank adapters in LLMs

[Hu et al. 2022] [Bershtatsky et al. 2024]

why tensor factorizations? (4/4)

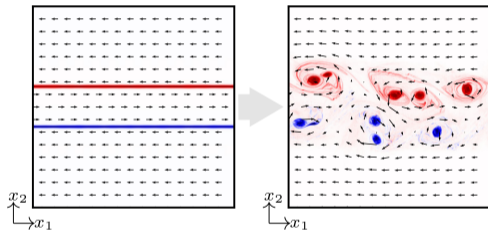
tensors 4 physics



**Fluid velocity vectors computed
in exponentially many points...**

why tensor factorizations? (4/4)

tensors 4 physics



Fluid velocity vectors computed in exponentially many points...

$$U_i \approx \begin{matrix} U_i^1 & U_i^2 & U_i^3 \\ \bullet & \bullet & \bullet \\ \omega_1 & \omega_2 & \omega_3 \end{matrix}$$

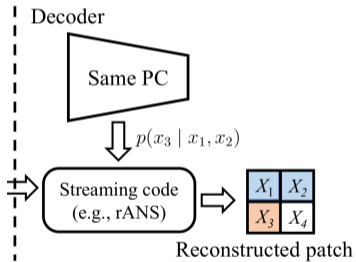
The diagram shows a tensor factorization. On the left, a blue horizontal bar represents a tensor U_i with three legs labeled ω_1 , ω_2 , and ω_3 . This is followed by an approximation symbol \approx . On the right, a chain of three tensors is shown: U_i^1 (red dot), U_i^2 (green dot), and U_i^3 (blue dot). The red dot has a leg labeled ω_1 . The green dot has a leg labeled ω_2 . The blue dot has a leg labeled ω_3 . The red and green dots are connected by a horizontal line labeled α_1 , and the green and blue dots are connected by a horizontal line labeled α_2 .

...by factorizing them into chains of low-rank tensors

[Gourianov et al. 2022] [Hölscher et al. 2025]

why circuits? (1/3)

efficient probabilistic inference

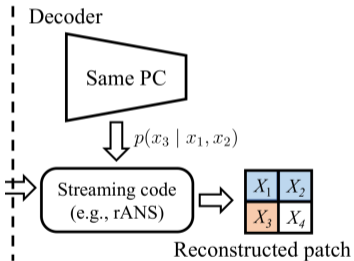


Fast lossless (de)compression

[Liu, Mandt, and Van den Broeck 2022]

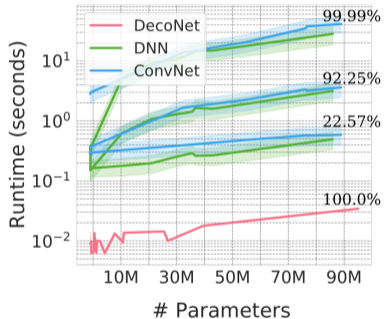
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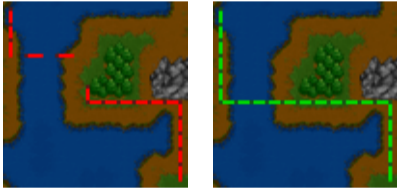


Efficient robustness to adversarial attacks

[Subramani et al. 2021]

why circuits? (2/3)

they enable neuro-symbolic AI

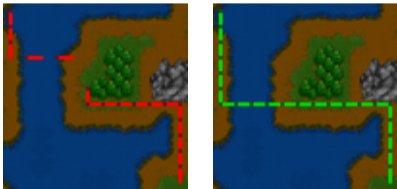


**Constrained multi-label
prediction (w/ guarantees)**

[Ahmed et al. 2022]

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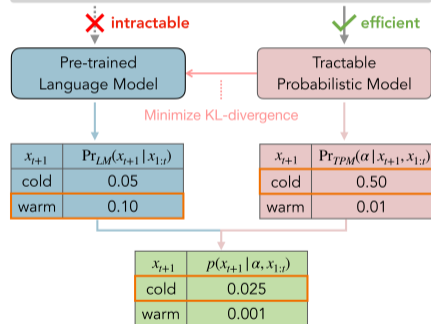


Constrained multi-label prediction (w/ guarantees)

[Ahmed et al. 2022]

Lexical Constraint α : sentence contains keyword "winter"

Constrained Generation: $\Pr(x_{t+1} | \alpha, x_{1:t} = \text{"the weather is"})$

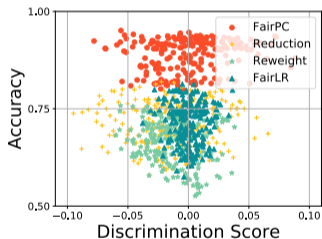


Constrained text generation

[Zhang et al. 2023]

why circuits? (3/3)

they are reliable and interpretable

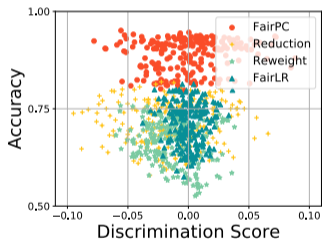


Encode group fairness

[Choi, Dang, and Van den Broeck 2020]

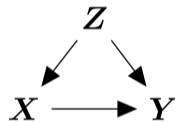
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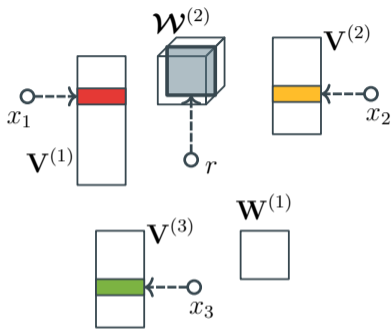
[Choi, Dang, and Van den Broeck 2020]



$$\sum_{\mathbf{Z}} p(\mathbf{Z})p(\mathbf{Y} | \mathbf{X}, \mathbf{Z})$$

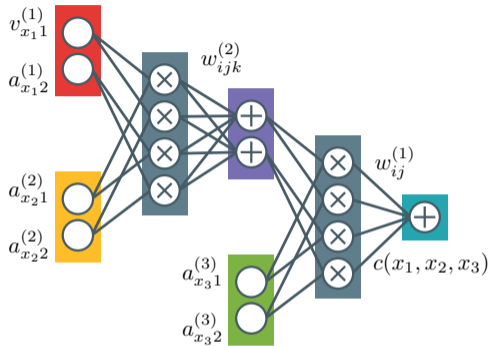
Tractable causal inference

[Wang and Kwiatkowska 2023]



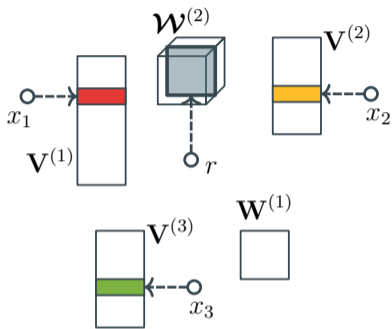
tensor factorizations

*tensor compression, graph data
physics-inspired AI, speed up LLMs...*

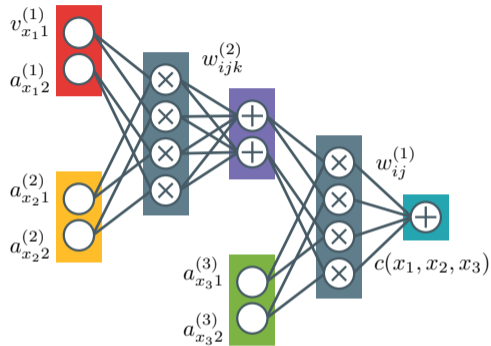


circuits

*property-driven fast inference
neuro-symbolic, trustworthy AI...*



tensor factorizations



circuits

two faces of the same coin...!

outline

1 connecting *tensor factorizations* and *circuits*

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- 2 a *unifying pipeline* to build factorizations & circuits

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- 3 a *property-driven* approach to inference & reasoning

outline

- 1 connecting *tensor factorizations* and *circuits*
- 2 a *unifying pipeline* to build factorizations & circuits
- 3 a *property-driven* approach to inference & reasoning
- 4 *expressiveness analysis*: known and new results

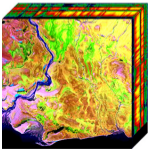
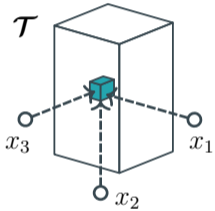
outline

1 connecting *tensor factorizations* and *circuits*

tl;dr

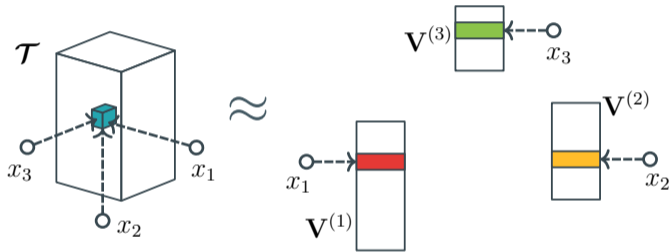
***“Understand when and how
a tensor factorization can be
exactly encoded as a circuit representation”***

Canonical polyadic (CP)



(N. M. Short)

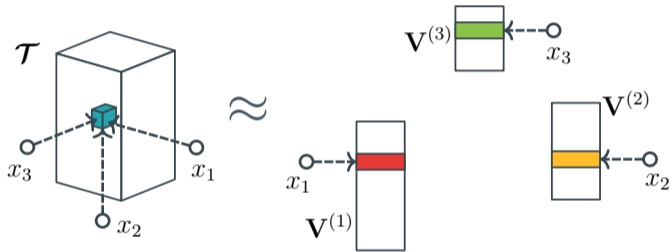
Canonical polyadic (CP)



$$t_{x_1 x_2 x_3} \approx \sum_{r=1}^R v_{x_1 r}^{(1)} v_{x_2 r}^{(2)} v_{x_3 r}^{(3)}$$

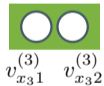
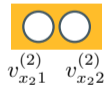
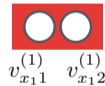
$$\mathbf{V}^{(1)} \in \mathbb{R}^{I_1 \times R}, \mathbf{V}^{(2)} \in \mathbb{R}^{I_2 \times R}, \mathbf{V}^{(3)} \in \mathbb{R}^{I_3 \times R}$$

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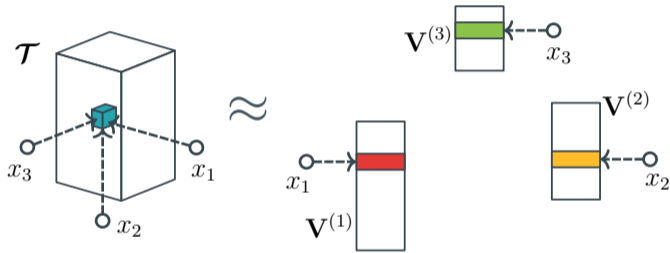


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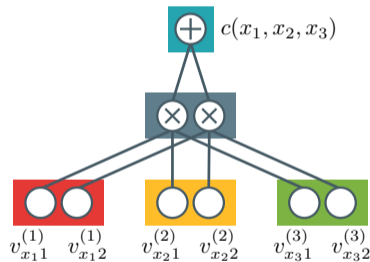


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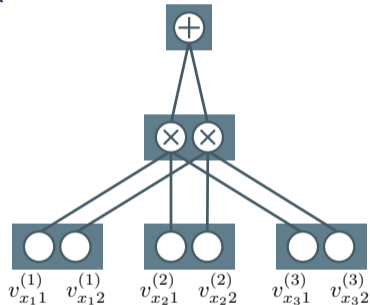
how to evaluate a circuit?

a circuit computes a tensor entry at some index

$$\mathbf{V}^{(1)} = \begin{bmatrix} 0.1 & 1.2 \\ 3.5 & -0.2 \\ -0.1 & 0.2 \end{bmatrix}$$

$$\mathbf{V}^{(2)} = \begin{bmatrix} 2.5 & 0.0 \\ -3.4 & -0.5 \\ -0.1 & 2.2 \end{bmatrix}$$

$$\mathbf{V}^{(3)} = \begin{bmatrix} -2.3 & 1.0 \\ 0.8 & -2.4 \\ 0.7 & 1.5 \end{bmatrix}$$

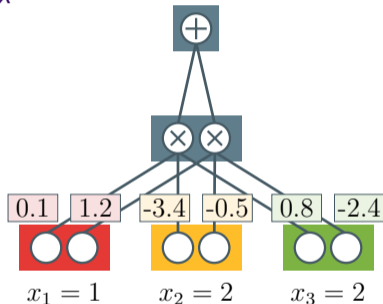


$$t_{x_1 x_2 x_3} \approx c(x_1, x_2, x_3) = \sum_{r=1}^R v_{x_1 r}^{(1)} v_{x_2 r}^{(2)} v_{x_3 r}^{(3)}$$

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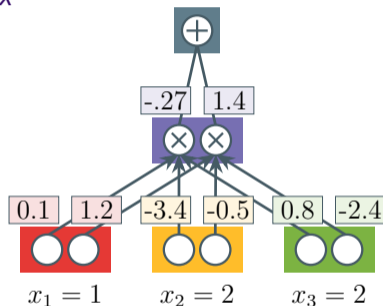


$$t_{122} \approx c(1, 2, 2) = \sum_{r=1}^R v_{1r}^{(1)} v_{2r}^{(2)} v_{2r}^{(3)}$$

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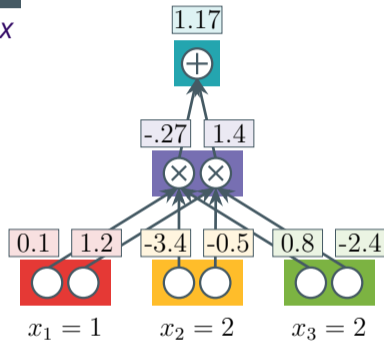


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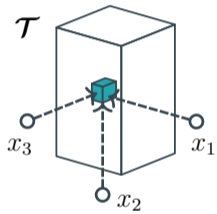
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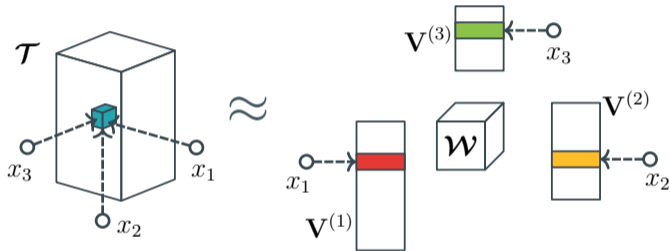


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Tucker



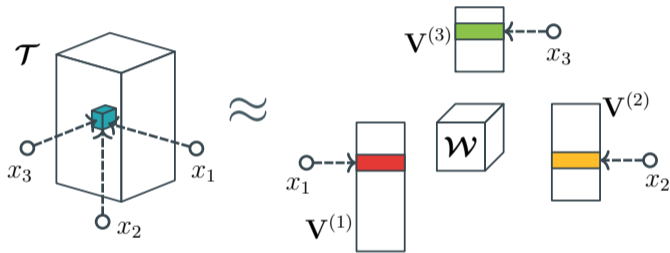
Tucker



$$t_{x_1 x_2 x_3} \approx \sum_{r_1=1}^{R_1} \sum_{r_2=1}^{R_2} \sum_{r_3=1}^{R_3} w_{r_1 r_2 r_3} v_{x_1 r_1}^{(1)} v_{x_2 r_2}^{(2)} v_{x_3 r_3}^{(3)}$$

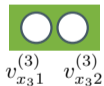
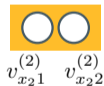
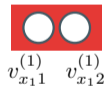
$$\mathcal{W} \in \mathbb{R}^{R_1 \times R_2 \times R_3}$$

Tucker

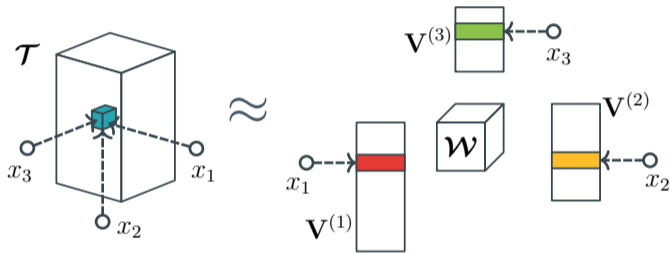


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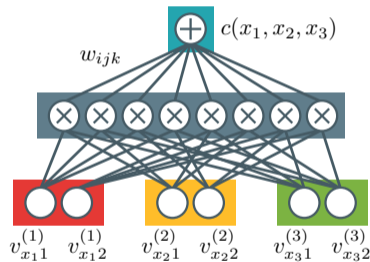


Tucker



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The layer-wise circuit definition

A collection of input units is a circuit layer $\ell(\mathbf{x}) \in \mathbb{R}^K$



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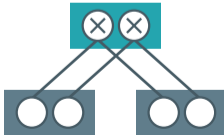


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A collection of input units is a circuit layer $\ell(\mathbf{x}) \in \mathbb{R}^K$

The product of two layers is a layer

$$\ell(\mathbf{x}) = \ell_i(\mathbf{x}) \odot \ell_{ii}(\mathbf{x}) \quad (\text{Hadamard})$$



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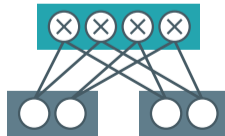
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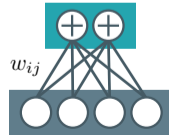
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A linear projection of a layer is a layer

$$\ell(\mathbf{x}) = \mathbf{W}\ell_i(\mathbf{x})$$



The layer-wise circuit definition

A collection of input units is a circuit layer $\ell(\mathbf{x}) \in \mathbb{R}^K$

The product of two layers is a layer

$$\ell(\mathbf{x}) = \ell_i(\mathbf{x}) \odot \ell_{ii}(\mathbf{x}) \quad (\text{Hadamard})$$

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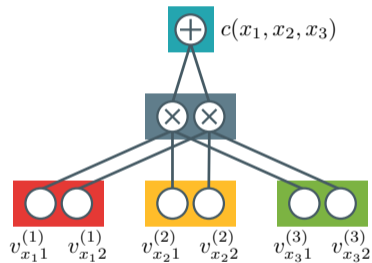
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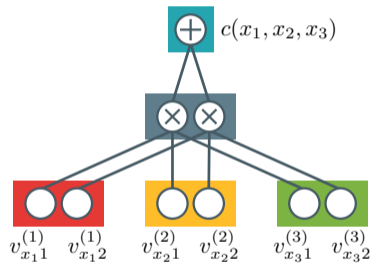


The layer-wise circuit definition

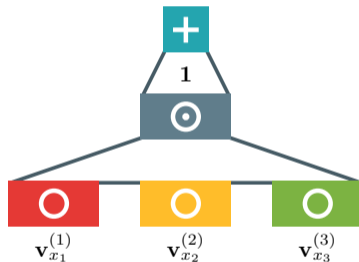


CP circuit

The layer-wise circuit definition

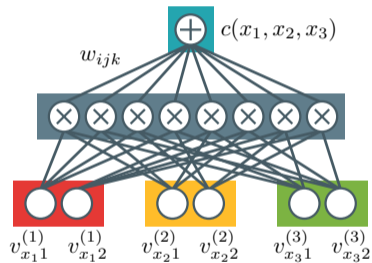


CP circuit



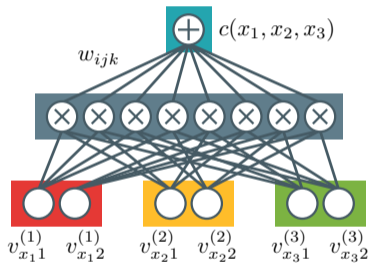
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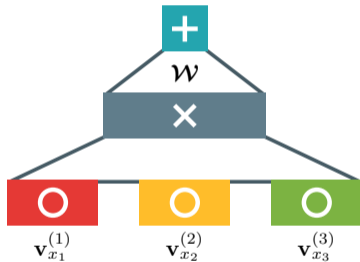


Tucker circuit

The layer-wise circuit definition



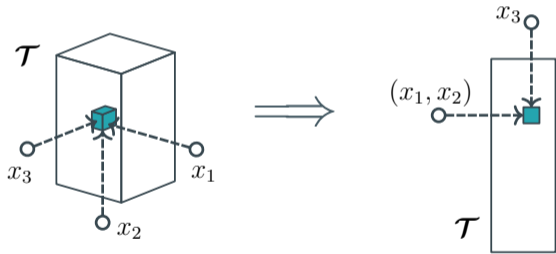
Tucker circuit



Tucker circuit

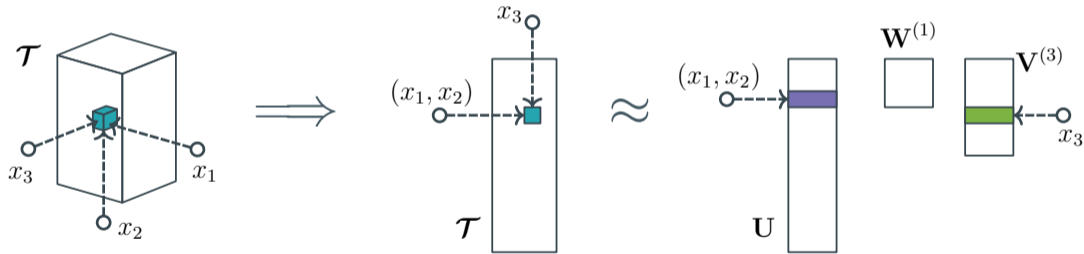
going deeper: hierarchical Tucker

level-one factorization



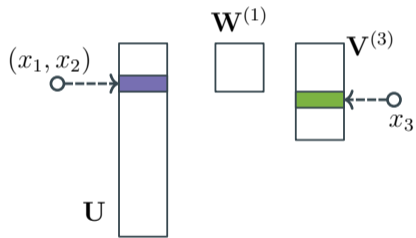
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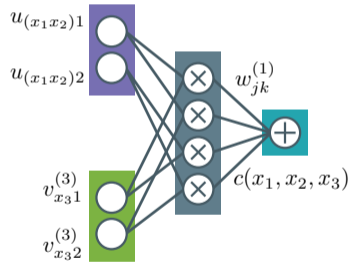
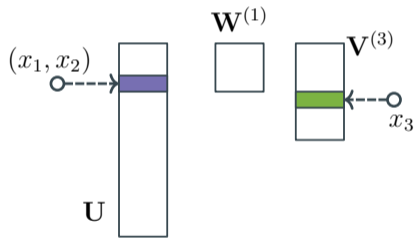
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level-one factorization as a circuit



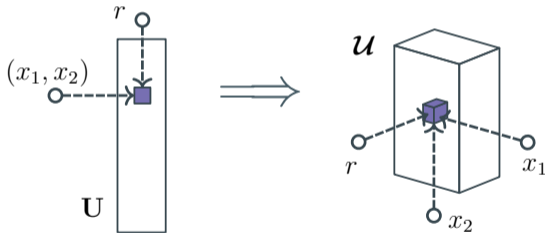
going deeper: hierarchical Tucker

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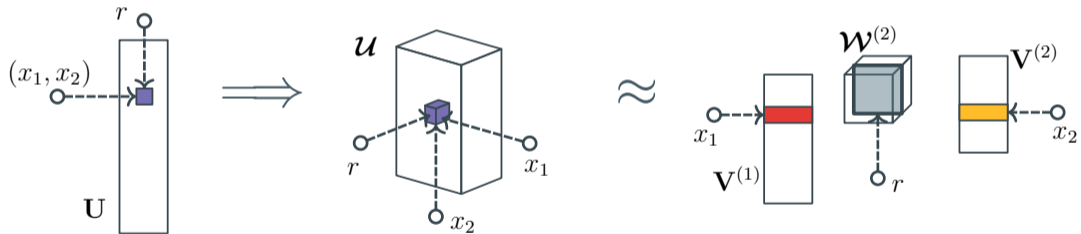
going deeper: hierarchical Tucker

level-two factorization



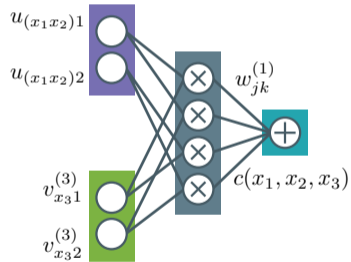
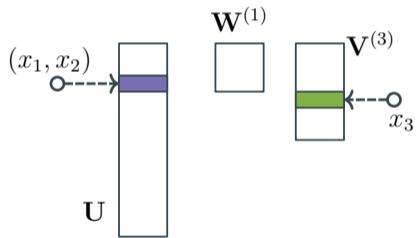
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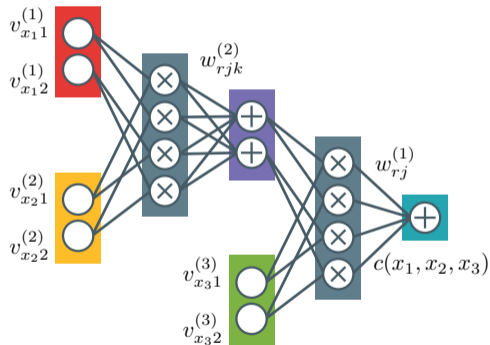
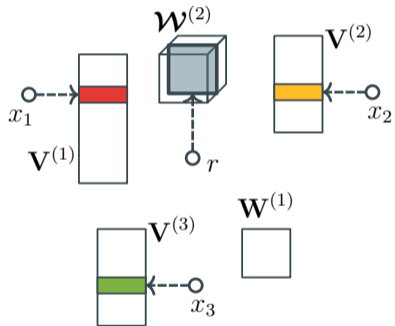
going deeper: hierarchical Tucker

nested factorizations are deep circuits



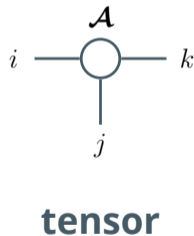
going deeper: hierarchical Tucker

nested factorizations are deep circuits



Tensor networks

the Penrose graphical notation



Tensor networks

matrix factorization & contraction



$$\mathbf{A} \in \mathbb{R}^{N \times R}$$



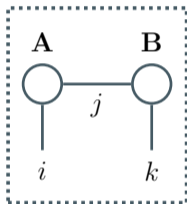
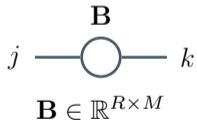
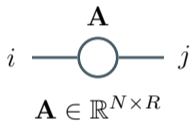
$$\mathbf{B} \in \mathbb{R}^{R \times M}$$

Biamonte and Bergholm, "Tensor Networks in a Nutshell", 2017

Orús, "A Practical Introduction to Tensor Networks: Matrix Product States and Projected Entangled Pair States", 2013

Tensor networks

matrix factorization & contraction



$$\mathbf{C} = \mathbf{A}\mathbf{B} \in \mathbb{R}^{N \times M}$$

$$c_{ik} = \sum_{j=1}^R a_{ij} b_{jk}$$

=

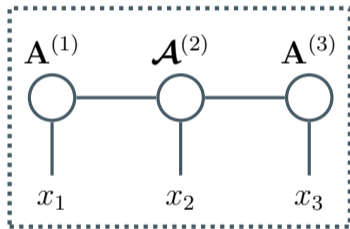


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Tensor trains are circuits

also called matrix-product states

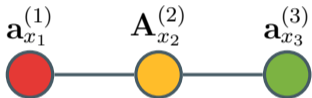


$$\mathcal{T} \in \mathbb{R}^{I_1 \times I_2 \times I_3}$$

$$t_{x_1 x_2 x_3} = \sum_{r_1=1}^R \sum_{r_2=1}^R a_{x_1 r_1}^{(1)} a_{r_1 x_2 r_2}^{(2)} a_{r_2 x_3}^{(3)}$$

Tensor trains are circuits

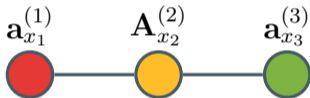
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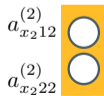
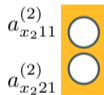
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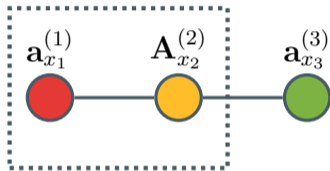


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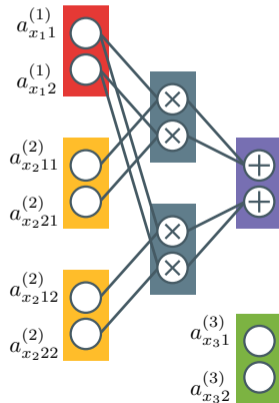


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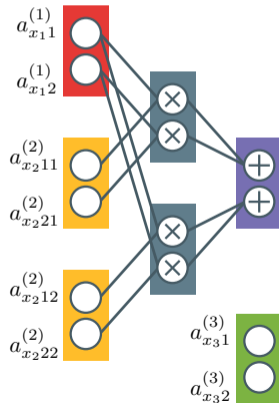


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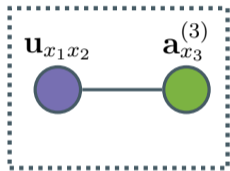


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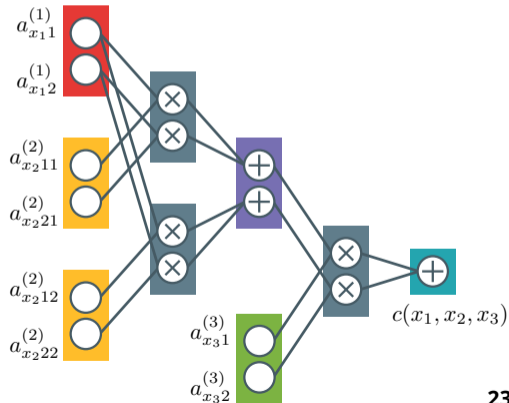


Tensor trains are circuits

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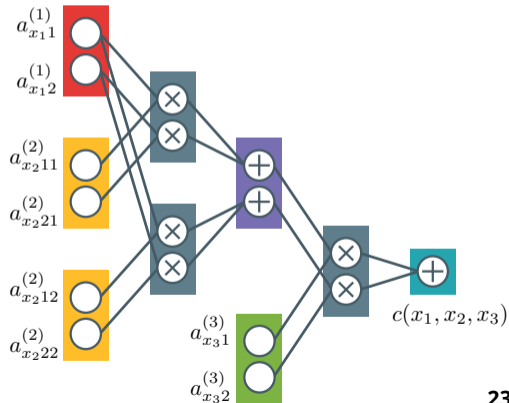
Tensor trains are circuits

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$t_{x_1 x_2 x_3}$



$t_{x_1 x_2 x_3}$



Many tensor factorizations are circuits

CP

RESICAL

Tucker

Hierarchical Tucker

Tensor train

Matrix-product state

Hierarchical Tucker

Tree tensor network

Complex

⋮

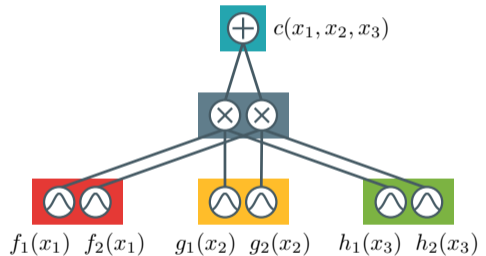
“What do we gain from circuits?”

More input functions with circuits

Input unit functions $f(x)$ compute:

— an entry of matrix (or tensor):

$$f(x) = v_{xr} \quad (\text{embedding layer})$$



More input functions with circuits

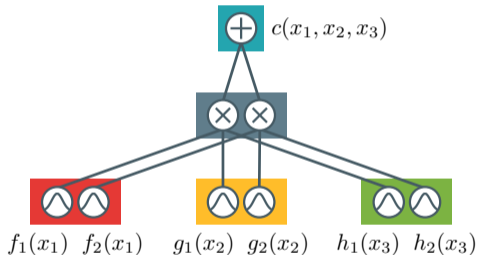
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$$f(x) = \text{Binomial}(x; n, p) \quad \text{(more compact!)}$$



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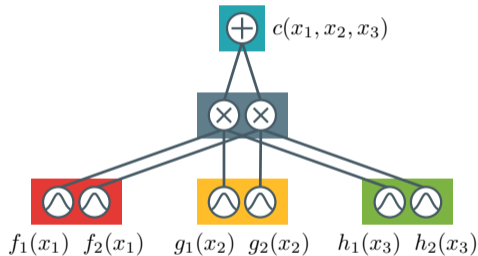
$$f(x) = \text{Binomial}(x; n, p) \quad \text{(more compact!)}$$

— continuous functions:

$$f(x) = a_0 + a_1x + \dots + a_nx^n$$

$$f(x) = \text{Normal}(x; \mu, \sigma^2)$$

⇒ **infinite-dimensional tensors (or functions)**



Townsend and Trefethen, "Continuous analogues of matrix factorizations", 2015

Novikov, Panov, and Oseledets, "Tensor-train density estimation", 2021

Probabilistic circuits (PCs)

PC == a circuit c encoding a non-negative function

$$\forall \mathbf{x} \in \text{dom}(\mathbf{X}): c(\mathbf{x}) = c(x_1, \dots, x_n) \geq 0$$

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$$p(\mathbf{x}) = \frac{1}{Z} c(\mathbf{x}),$$

where $Z = \sum_{\mathbf{x}} c(\mathbf{x})$ (PMF) or $Z = \int c(\mathbf{x}) d\mathbf{x}$ (PDF)

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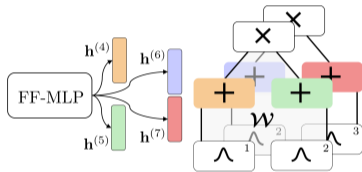
Non-negative sum parameters \wedge non-negative input functions

\implies a circuit is a **PC**

Cichocki and Phan, "Fast Local Algorithms for Large Scale Nonnegative Matrix and Tensor Factorizations", 2009

How to parameterize circuits?

(i.e., the weights of sums and input functions)



Functions: neural networks

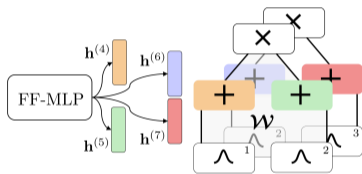
Gala et al., "Scaling Continuous Latent Variable Models as Probabilistic Integral Circuits", 2024

Shao et al., "Conditional sum-product networks: Imposing structure on deep probabilistic architectures", 2020

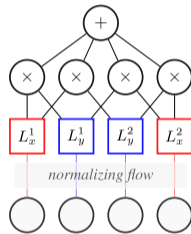
Sidheekh, Kersting, and Natarajan, "Probabilistic Flow Circuits: Towards Unified Deep Models for Tractable Probabilistic Inference", 2023

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Deep generative models

Gala et al., "Scaling Continuous Latent Variable Models as Probabilistic Integral Circuits", 2024

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learning & reasoning with circuits in pytorch

`https://github.com/april-tools/circuit`



```
1 from cirq.symbolic.layers import (  
2     EmbeddingLayer, SumLayer, KroneckerLayer,  
3     Scope  
4 )  
5  
6 # Tensor shape and rank  
7 shape = (3, 1280, 720)  
8 rank = 42  
9  
10 # Construct the Tucker factorization layers  
11 v1 = EmbeddingLayer(Scope([0]), rank, num_states=shape[0])  
12 v2 = EmbeddingLayer(Scope([1]), rank, num_states=shape[1])  
13 v3 = EmbeddingLayer(Scope([2]), rank, num_states=shape[2])  
14 kron = KroneckerLayer(rank, arity=3)  
15 tucker = SumLayer(rank ** 3, num_output_units=1)
```

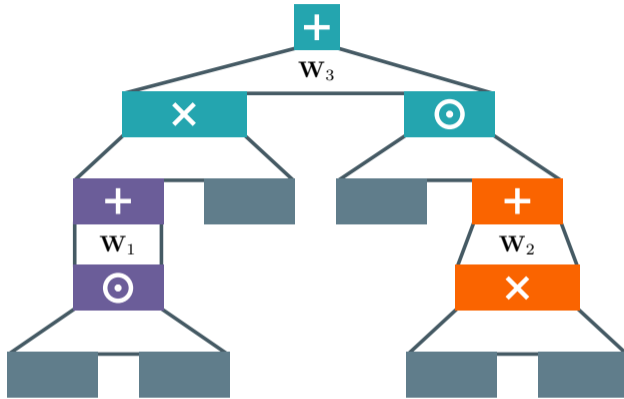


```
17  
18 # Construct the Tucker circuit
```

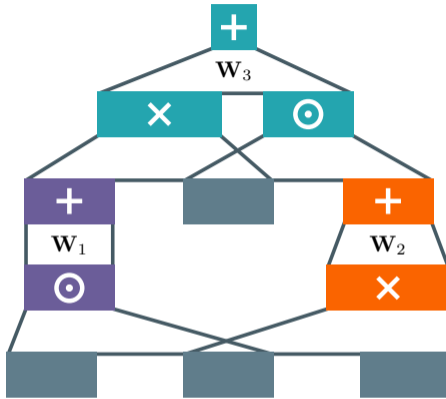
```
19 circuit = Circuit(  
20     layers=[v1, v2, v3, kron, tucker],  
21     in_layers={ # The layers input connections  
22         kron: [v1, v2, v3],  
23         tucker: [kron]  
24     },  
25     outputs=[tucker]  
26 )
```

```
1 # Compile the circuit to PyTorch code
2 from cirkit.pipeline import compile
3 pth_circuit = compile(circuit)
4
5 print(pth_circuit) # Tucker factorization
6 # TorchCircuit(
7 #   (0): TorchEmbeddingLayer(...)
8 #   (1): TorchEmbeddingLayer(...)
9 #   (2): TorchEmbeddingLayer(...)
10 #   (3): TorchKroneckerLayer(...)
11 #   (4): TorchSumLayer(...)
12 #)
13
14 # Compute one entry of the encoded tensor
15 x = torch.tensor([[1, 500, 300]])
16 t_x = pth_circuit(x)
```

```
1 from cirkit.templates import tensor_factorizations
2
3 shape = (3, 1280, 720)
4
5 # CP factorization
6 circuit = tensor_factorizations.cp(shape, rank=42)
7
8 # Tucker factorization
9 circuit = tensor_factorizations.tucker(shape, rank=42)
10
11 # Tensor-train / matrix-product state
12 circuit = tensor_factorizations.tensor_train(shape, rank=42)
```



Stack layers to build a deep factorization!



Save computation by sharing sub-factorizations!



```
1 from cirq.symbolic.layers import (  
2     EmbeddingLayer, SumLayer, KroneckerLayer,  
3     HadamardLayer, Scope)  
4  
5 # Tensor shape and ranks  
6 shape = (17, 3, 1280, 720)  
7 rank1, rank2 = 2, 4  
8  
9 # Construct the layers  
10 v1 = EmbeddingLayer(Scope([0]), rank1, num_states=shape[0])  
11 v2 = EmbeddingLayer(Scope([1]), rank1, num_states=shape[1])  
12 v3 = EmbeddingLayer(Scope([2]), rank1, num_states=shape[2])  
13 v4 = EmbeddingLayer(Scope([3]), rank1, num_states=shape[3])  
14 kron1 = KroneckerLayer(rank1, arity=2)  
15 kron2 = KroneckerLayer(rank1, arity=2)  
16 hada1 = HadamardLayer(rank1, arity=2)
```

```
17 hada2 = HadamardLayer(rank1, arity=2)
18 sum1 = SumLayer(rank1, num_output_units=rank1)
19 sum2 = SumLayer(rank1 ** 2, num_output_units=rank1)
20 sum3 = SumLayer(rank1 + rank2, num_output_units=1, arity=2)
21
22 # Construct the "Frankenstein" circuit
23 circuit = Circuit(
24     layers=[v1, v2, v3, v4, kron1, kron2, ...],
25     in_layers={ # The layers input connections
26         hada1: [v1, v2],
27         kron1: [v2, v3],
28         sum1: [hada1],
29         sum2: [kron1],
30         kron2: [sum1, v4],
31         ...
```

Takeaways

- 1 circuits *unify* many (deep) tensor factorizations

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2 A *Lego block* approach to tensor factorizations
(*different parameterizations*)

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- 1 circuits **unify** many (deep) tensor factorizations
- 2 A **Lego block** approach to tensor factorizations
(different parameterizations)
- 3 build **new** tensor factorizations by connecting layers
*(easy to do within the **cirkit** library)*

Takeaways

Questions?

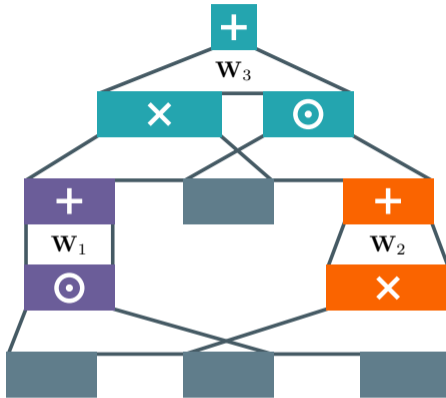
- 1 circuits **unify** many (deep) tensor factorizations
- 2 A **Lego block** approach to tensor factorizations
(*different parameterizations*)
- 3 build **new** tensor factorizations by connecting layers
(*easy to do within the **cirkit** library*)

outline

- 1 connecting *tensor factorizations* and *circuits*
- 2 a *unifying pipeline* to build factorizations & circuits

tl;dr

***“Understand when and how
we can build a deep circuit
that is a deep factorization”***



...but do we always get a tensor factorization?

Multilinear forms

tensor factorizations are typically multilinear

linear in each univariate basis/input function

$$\sum_i \alpha_i \prod_{j=1}^d f_{i,j}(x_j) \left\{ \begin{array}{l} t_{x_1 \dots x_n} = \sum_{r=1}^R \prod_{j=1}^d v_{x_j r}^{(j)} \quad (\text{CP}) \\ t_{x_1 \dots x_n} = \sum_{r_1, \dots, r_d=1}^{R_1, \dots, R_d} w_{r_1 \dots r_d} \prod_{j=1}^d v_{x_j r_j}^{(j)} \quad (\text{Tucker}) \end{array} \right.$$

Vasilescu and Terzopoulos, "Multilinear Image Analysis for Facial Recognition", 2002
Kolda, Multilinear operators for higher-order decompositions, 2006

“How to enforce multilinearity in deep circuits?”

Structural properties

smoothness

decomposability

compatibility

Structural properties

property A

property B

property C

Structural properties

smoothness

decomposability

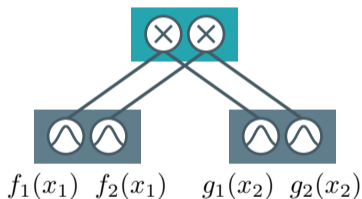
property C

smoothness \wedge *decomposability*

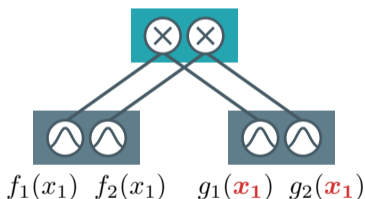
\implies **multilinearity**

Multilinearity in circuits

the inputs of product units are defined over disjoint sets of variables



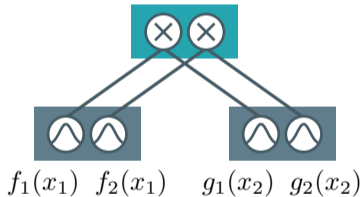
✓ **multilinear**



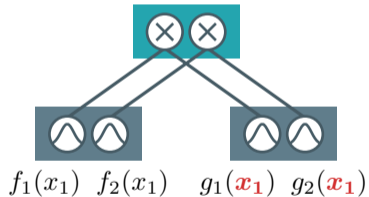
✗ **not multilinear**

Multilinearity in circuits

the inputs of product units are defined over disjoint sets of variables



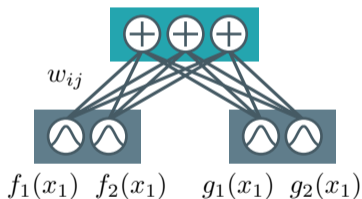
decomposable circuit



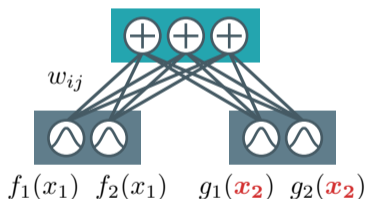
non-decomposable circuit

Multilinearity in circuits

the inputs of sum units are defined over the same variables



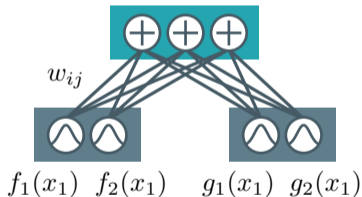
✓ multilinear



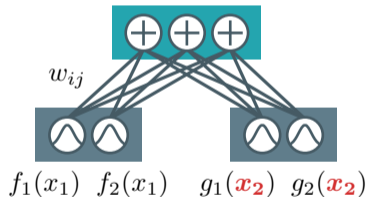
✗ not multilinear

Multilinearity in circuits

the inputs of sum units are defined over the same variables



smooth circuit



non-smooth circuit

Structural properties

smoothness

decomposability

property C

smoothness \wedge *decomposability*

\implies **multilinearity**

Structural properties

smoothness

decomposability

property C

tractable computation of **arbitrary integrals**
in probabilistic circuits

$$p(\mathbf{y}) = \int p(\mathbf{y}, \mathbf{z}) d\mathbf{z}, \quad \forall \mathbf{Y} \subseteq \mathbf{X}, \quad \mathbf{Z} = \mathbf{X} \setminus \mathbf{Y}$$

\implies tractable partition function

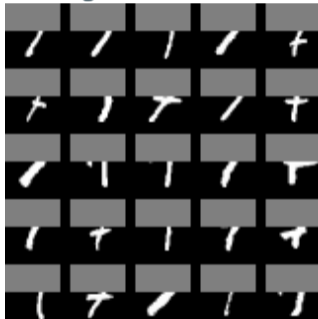
\implies also any conditional is tractable

tractable marginals on PCs

Original

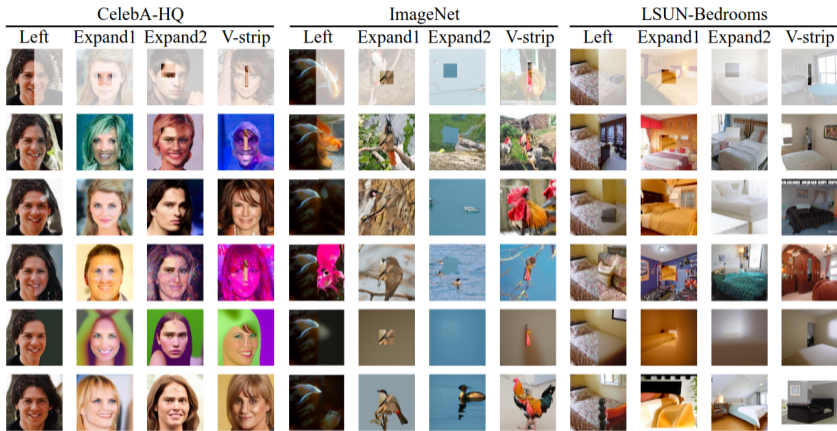


Missing



Conditional sample





smooth + **decomposable** circuits = ...

compute arbitrary summations (or integrals)
⇒ linear time in circuit size!

E.g., partition function $\sum_{\mathbf{x}} c(\mathbf{x})$

or $\int c(\mathbf{x}) d\mathbf{x}$ in the continuous case

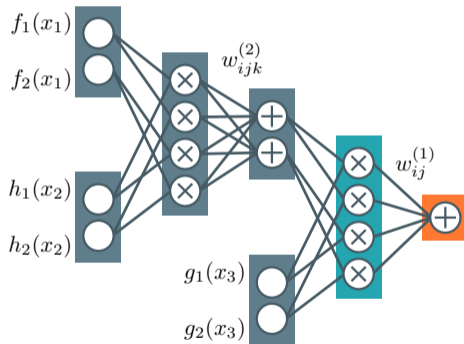
smooth + **decomposable** circuits = ...

$$\text{If } c(\mathbf{x}) = \sum_i w_i c_i(\mathbf{x})$$

(smoothness):

$$\begin{aligned} \int c(\mathbf{x}) \, d\mathbf{x} &= \int \sum_i w_i c_i(\mathbf{x}) \, d\mathbf{x} \\ &= \sum_i w_i \int c_i(\mathbf{x}) \, d\mathbf{x} \end{aligned}$$

⇒ integrals are “pushed down” to the inputs



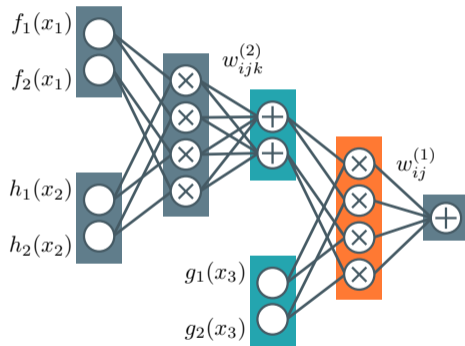
smooth + **decomposable** circuits = ...

$$\text{If } c(\mathbf{x}) = c_1(\mathbf{y}) c_2(\mathbf{z})$$

(decomposability):

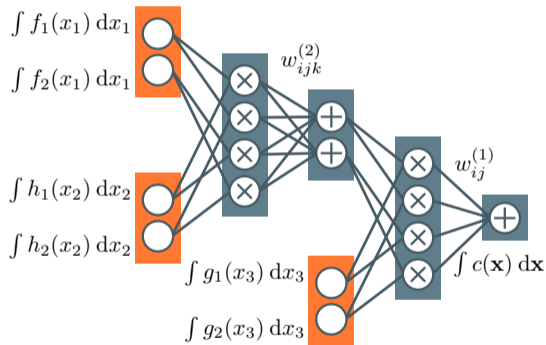
$$\begin{aligned} \int c(\mathbf{x}) \, d\mathbf{x} &= \int \int c_1(\mathbf{y}) c_2(\mathbf{z}) \, d\mathbf{y} d\mathbf{z} \\ &= \left(\int c_1(\mathbf{y}) \, d\mathbf{y} \right) \left(\int c_2(\mathbf{z}) \, d\mathbf{z} \right) \end{aligned}$$

\Rightarrow integrals “decompose” into easier ones



smooth + **decomposable** circuits = ...

Integrate simple input functions $f(x)$
⇒ Gaussians, polynomials, splines, ...



“How to build smooth & decomposable circuits?”

“Can we re-use known tensor factorization methods?”

A zoo of probabilistic circuits...

PC ARCHITECTURE

Poon&Domingos (Poon & Domingos, 2011)

RAT-SPN (Peharz et al., 2020c)

EiNet (Peharz et al., 2020a)

HCLT (Liu & Van den Broeck, 2021b)

HMM/MPS $_{\mathbb{R} \geq 0}$ (Glasser et al., 2019)

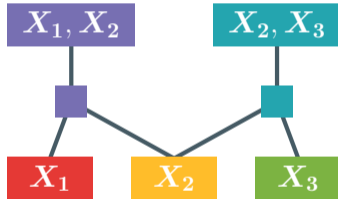
BM (Han et al., 2018)

TTDE (Novikov et al., 2021)

NPC² (Loconte et al., 2024)

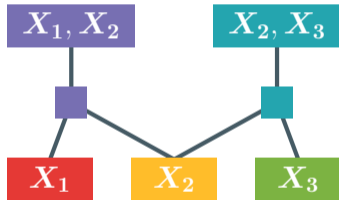
TTN (Cheng et al., 2019)

Building smooth & decomposable circuits

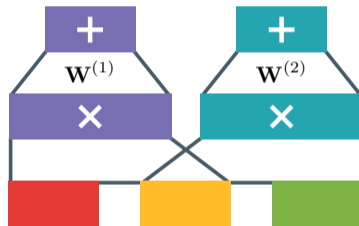


1) choose a template

Building smooth & decomposable circuits



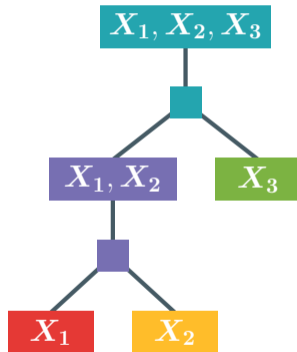
1) choose a template



2) pick a layer to parameterize the chosen template

Region graphs

A bipartite graph to build smooth and decomposable circuits:



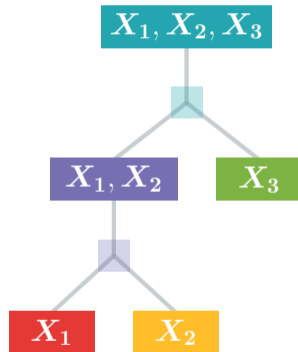
Dennis and Ventura, "Learning the architecture of sum-product networks using clustering on variables", 2012

Region graphs

A bipartite graph to build smooth and decomposable circuits:

Region node:

set of variables (or dimensions)



Region graphs

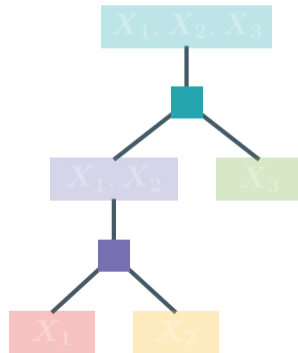
A bipartite graph to build smooth and decomposable circuits:

Region node:

set of variables (or dimensions)

Partition node:

decomposition of a region



Region graphs

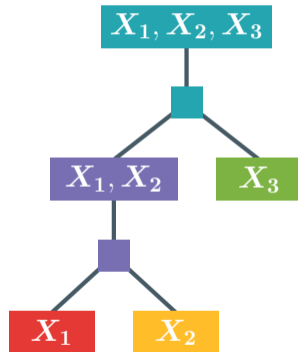
A bipartite graph to build smooth and decomposable circuits:

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Region graphs

A bipartite graph to build smooth and decomposable circuits:

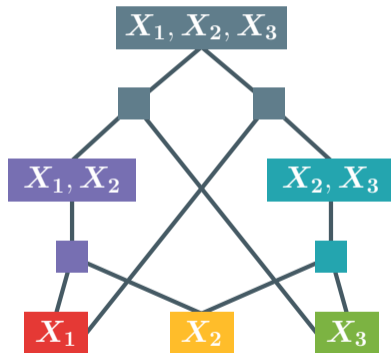
Region node:

set of variables (or dimensions)

Partition node:

decomposition of a region

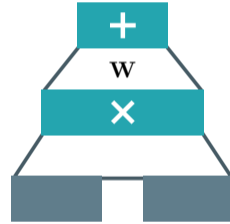
⇒ generalizes mode cluster trees



Circuit sum-product layers as factorizations

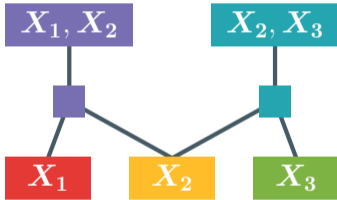


CP layer



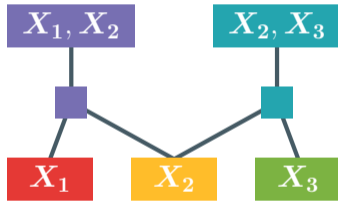
Tucker layer

From region graphs to circuits

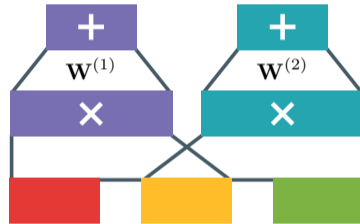


1) choose a region graph

From region graphs to circuits

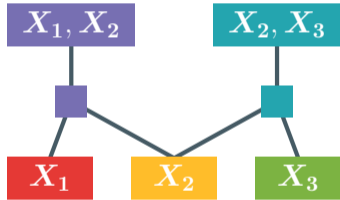


1) choose a region graph

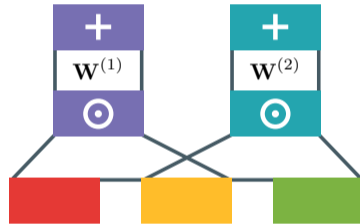


2) pick layer and number of units
(e.g., Tucker layer)

From region graphs to circuits



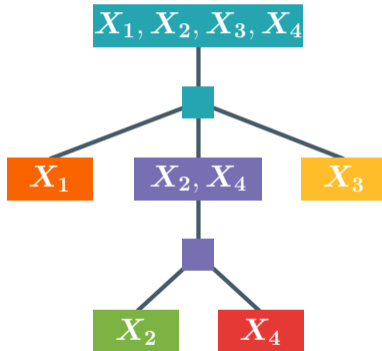
1) choose a region graph



2) pick layer and number of units
(e.g., CP layer)

Which region graph?

learned or randomized trees

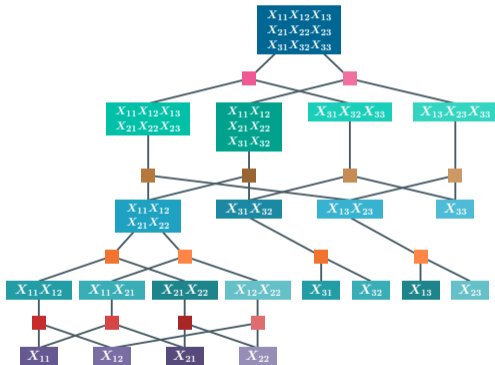


Peharz et al., "Random Sum-Product Networks: A Simple and Effective Approach to Probabilistic Deep Learning", 2020

Liu and Broeck, "Tractable Regularization of Probabilistic Circuits", 2021

Which region graph?

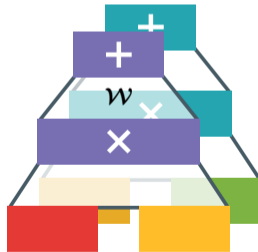
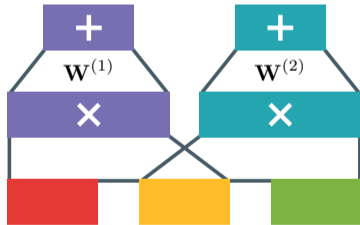
image-tailored graphs



folding

to speed-up inference

parallelize layers that can be evaluated independently



A unifying circuit construction pipeline

PC ARCHITECTURE	REGION GRAPH	SUM-PRODUCT LAYER	FOLD
Poon&Domingos (Poon & Domingos, 2011)	PD	CP^T	✗
RAT-SPN (Peharz et al., 2020c)	RND	Tucker	✗
EiNet (Peharz et al., 2020a)	{ RND, PD }	Tucker	✓
HCLT (Liu & Van den Broeck, 2021b)	CL	CP^T	✓
HMM/MPS _{$\mathbb{R}_{\geq 0}$} (Glasser et al., 2019)	LT	CP^T	✗
BM (Han et al., 2018)	LT	CP^T	✗
TTDE (Novikov et al., 2021)	LT	CP^T	✗
NPC ² (Loconte et al., 2024)	{ LT, RND }	{ CP^T , Tucker }	✓
TTN (Cheng et al., 2019)	QT-2	Tucker	✗
Mix & Match (our pipeline)	{ RND, PD, LT, CL, QG, QT-2, QT-4 }	{ Tucker, CP, CP^T } \cup { CP^S , CP^{XS} FOLD ✓ }	{ ✗, ✓ }

enrich the pipeline

with new layers

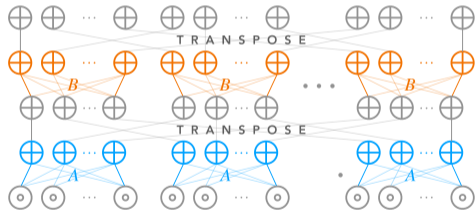
Monarch matrix factorization:

$$\mathbf{W}_M = \mathbf{P}_L \mathbf{L} \mathbf{P}_R \mathbf{R}$$

Monarch circuit layer:

$$\ell(\mathbf{x}) = \mathbf{W}_M \ell_i(\mathbf{x})$$

More tensor factorizations
with new layers!





```
1 # Construct a region graph
2 from cirq.templates.region_graph import (
3     RandomBinaryTree) # or QuadGraph, LinearTree, ...
4 region_graph = RandomBinaryTree(num_variables=10)
5
6 # Build the circuit from the region graph
7 from cirq.symbolic.layers import EmbeddingLayer
8 circuit = region_graph.build_circuit(
9     sum_product='tucker', # or 'cp'
10    input_factory=EmbeddingLayer # or GaussianLayer, ...
11    num_sum_units=32,
12    num_input_units=32)
13
14 # Compile the circuit to PyTorch
15 from cirq.pipeline import compile
16 pth_circuit = compile(circuit)
```

Takeaways

- 1 **structural properties** for multilinearity

Takeaways

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- 2 **exact & efficient (*tractable*) summations**
(more properties & operations next!)

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- 1 **structural properties** for multilinearity
- 2 **exact & efficient (*tractable*) summations**
(more properties & operations next!)
- 3 **a *pipeline* to build circuits & tensor factorizations**
(different layers and graph structures)

Takeaways

Questions?

- 1 **structural properties** for multilinearity
- 2 **exact & efficient (*tractable*) summations**
(more properties & operations next!)
- 3 a **pipeline** to build circuits & tensor factorizations
(different layers and graph structures)

outline

- 1 connecting *tensor factorizations* and *circuits*
- 2 a *unifying pipeline* to build factorizations & circuits
- 3 a *property-driven* approach to inference & reasoning

tl;dr

***“Understand when and how
we can build a deep factorization
that guarantees tractable reasoning”***

reasoning about ML models



q₁

“What is the probability of a treatment for a patient with **unavailable records**?”



q₂

“How **fair** is the prediction with respect to protected attribute changes?”



q₃

“Can we certify no **adversarial examples** exist?”

Reasoning about ML models



q₁

$$\int p(\mathbf{x}_o, \mathbf{x}_m) d\mathbf{X}_m$$

(missing values)

q₂

$$\mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{x}_c | X_s=0)} [f_0(\mathbf{x}_c)] - \mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{x}_c | X_s=1)} [f_1(\mathbf{x}_c)]$$

(fairness)

q₃

$$\mathbb{E}_{\mathbf{e} \sim p_{\text{noise}}(\mathbf{E})} [f(\mathbf{x} + \mathbf{e})]$$

(adversarial robust.)

...in the language of probabilities

Reasoning about ML models



q₁

$$\int p(\mathbf{x}_o, \mathbf{x}_m) d\mathbf{X}_m$$

(missing values)

q₂

$$\mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{x}_c | X_s=0)} [f_0(\mathbf{x}_c)] - \mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{x}_c | X_s=1)} [f_1(\mathbf{x}_c)]$$

(fairness)

q₃

$$\mathbb{E}_{\mathbf{e} \sim p_{\text{noise}}(\mathbf{E})} [f(\mathbf{x} + \mathbf{e})]$$

(adversarial robust.)

hard to compute in general!

Reasoning about ML models



$$q_1 \int p(\mathbf{x}_o, \mathbf{x}_m) d\mathbf{x}_m$$

(missing values)

$$q_2 \mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{x}_c | X_s=0)} [f_0(\mathbf{x}_c)] - \mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{x}_c | X_s=1)} [f_1(\mathbf{x}_c)]$$

(fairness)

$$q_3 \mathbb{E}_{\mathbf{e} \sim p_{\text{noise}}(\mathbf{E})} [f(\mathbf{x} + \mathbf{e})]$$

(adversarial robust.)

it is crucial we compute them **exactly** and in **polytime!**

Which structural properties

for complex reasoning



q₁

$$\int p(\mathbf{x}_o, \mathbf{x}_m) d\mathbf{X}_m$$

(missing values)



q₂

$$\mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{x}_c | X_s=0)} [f_0(\mathbf{x}_c)] - \mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{x}_c | X_s=1)} [f_1(\mathbf{x}_c)]$$

(fairness)



q₃

$$\mathbb{E}_{\mathbf{e} \sim p_{\text{noise}}(\mathbf{E})} [f(\mathbf{x} + \mathbf{e})]$$

(adversarial robust.)

smooth + decomposable

Which structural properties

for complex reasoning



q₁

$$\int p(\mathbf{x}_o, \mathbf{x}_m) d\mathbf{X}_m$$

(missing values)

smooth + decomposable



q₂

$$\mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{x}_c | X_s=0)} [f_0(\mathbf{x}_c)] - \mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{x}_c | X_s=1)} [f_1(\mathbf{x}_c)]$$

(fairness)

???????



q₃

$$\mathbb{E}_{\mathbf{e} \sim p_{\text{noise}}(\mathbf{E})} [f(\mathbf{x} + \mathbf{e})]$$

(adversarial robust.)

???????

Which properties for expectations?

smoothness

decomposability

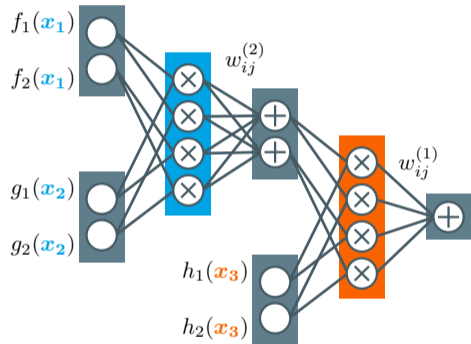
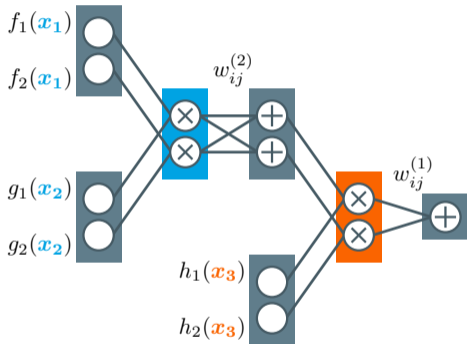
compatibility

Integrals involving two or more functions:
e.g., expectations

$$\mathbb{E}_{\mathbf{x} \sim p} f(\mathbf{x}) = \int p(\mathbf{x}) f(\mathbf{x}) d\mathbf{x}$$

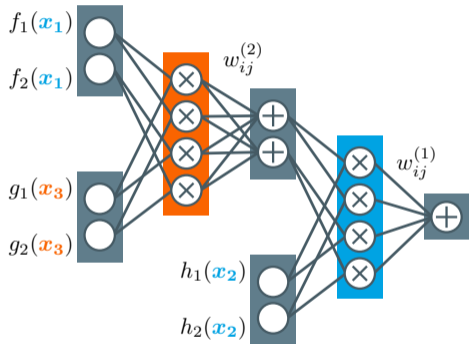
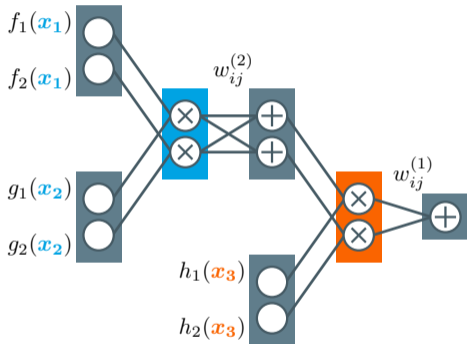
when both $p(\mathbf{x})$ and $f(\mathbf{x})$ are circuits

Compatibility



Compatible circuits

Compatibility



non-compatible circuits

Structural properties

smoothness

decomposability

compatibility

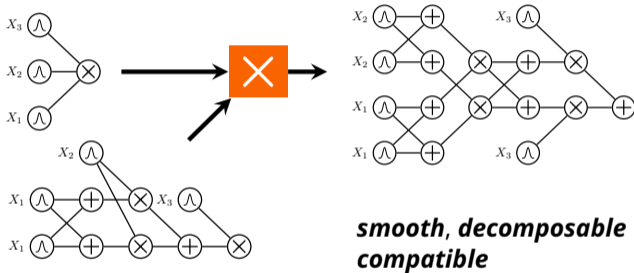
compatibility



smoothness \wedge decomposability

compatibility \Rightarrow tractable expectations

Tractable products



$$\text{compute } \mathbb{E}_{\mathbf{x} \sim p} f(\mathbf{x}) = \int p(\mathbf{x}) f(\mathbf{x}) \, d\mathbf{x} \text{ in } O(\|p\| \|f\|)$$



```
1 from cirq.symbolic.circuit import Circuit
2 from cirq.symbolic.functional import (
3     integrate, multiply)
```

```
4
5 # Circuits expectation  $\int p(x) f(x) dx$ 
```

```
6 def expectation(p: Circuit, f: Circuit) -> Circuit:
7     i = multiply(p, f)
8     return integrate(i)
```

```
9
10 # Squared loss  $\int [p(x)-q(x)]^2 dx = E_p[p] + E_q[q] - 2E_p[q]$ 
11 #           =  $\int p^2(x) dx + \int q^2(x) dx - 2\int p(x)q(x) dx$ 
```

```
12 def squared_loss(p: Circuit, q: Circuit) -> Circuit:
13     p2 = multiply(p, p)
14     q2 = multiply(q, q)
15     pq = multiply(p, q)
16     return integrate(p2) + integrate(q2) - 2 * integrate(pq)
```

Which structural properties

for complex reasoning



q₁

$$\int p(\mathbf{x}_o, \mathbf{x}_m) d\mathbf{X}_m$$

(missing values)

smooth + decomposability



q₂

$$\mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{x}_c | X_s=0)} [f_0(\mathbf{x}_c)] - \mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{x}_c | X_s=1)} [f_1(\mathbf{x}_c)]$$

(fairness)

compatibility



q₃

$$\mathbb{E}_{\mathbf{e} \sim p_{\text{noise}}(\mathbf{E})} [f(\mathbf{x} + \mathbf{e})]$$

(adversarial robust.)

compatibility

What if compatibility does not apply?

$$\mathbb{E}_{\mathbf{e} \sim p_{\text{noise}}(\mathbf{E})} [f(\mathbf{x} + \mathbf{e})]$$

p_{noise}

a circuit

f

not a circuit (e.g., neural net)



Q3

$\mathbb{E}_{\mathbf{e} \sim p_{\text{noise}}(\mathbf{E})} [f(\mathbf{x} + \mathbf{e})]$
(*adversarial robust.*)

How to approximate it by sampling?

wait...!

***“How can we sample
from a deep factorization
or tensor network?”***

approximate inference

e.g., via sampling

We can use **autoregressive inverse transform sampling**:

$$x_1 \sim p(x_1), \quad x_i \sim p(x_i | \mathbf{x}_{<i}) \quad \text{for } i \in \{2, \dots, d\}$$

\Rightarrow can be slow for large dimensions, requires **inverting the CDF**

approximate inference

e.g., via sampling

We can use **autoregressive inverse transform sampling**:

$$x_1 \sim p(x_1), \quad x_i \sim p(x_i | \mathbf{x}_{<i}) \quad \text{for } i \in \{2, \dots, d\}$$

\Rightarrow can be slow for large dimensions, requires **inverting the CDF**

can we do better?

approximate inference

e.g., via sampling

We can use **autoregressive inverse transform sampling**:

$$x_1 \sim p(x_1), \quad x_i \sim p(x_i | \mathbf{x}_{<i}) \quad \text{for } i \in \{2, \dots, d\}$$

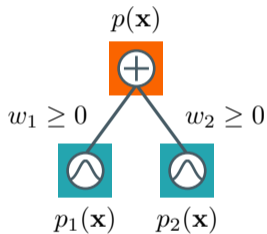
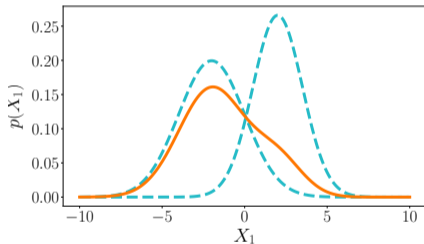
\Rightarrow can be slow for large dimensions, requires **inverting the CDF**

can we do better?

\Rightarrow yes, for non-negative factorizations/monotonic PCs

How to sample?

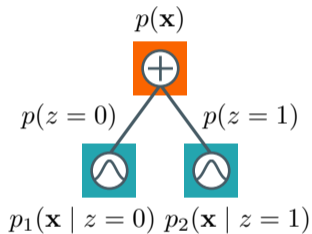
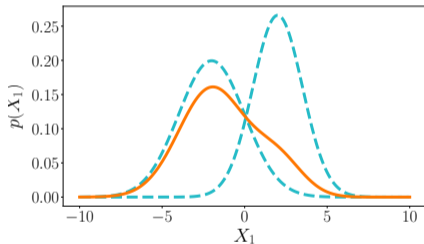
non-negative factorizations as latent-variable models



$$p(\mathbf{x}) = w_1 p_1(\mathbf{x}) + w_2 p_2(\mathbf{x})$$

How to sample?

non-negative factorizations as latent-variable models



$$p(\mathbf{x}) = p(z = 0) p_1(\mathbf{x} | z = 0) + p(z = 1) p_2(\mathbf{x} | z = 1)$$

Structural properties

smoothness

decomposability

compatibility

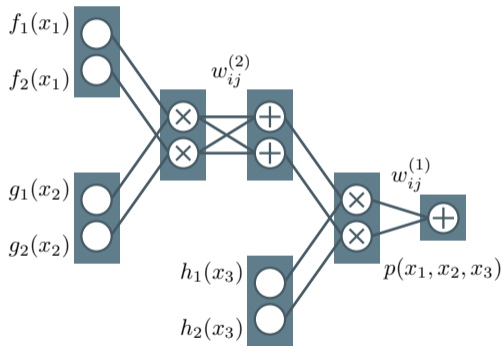
sampling in a single backward pass

draw $\mathbf{x} \sim p(\mathbf{X})$

\implies **exact** sampling method

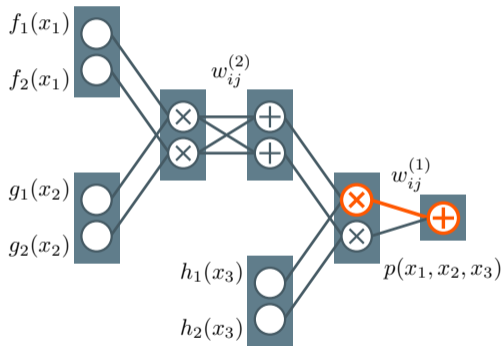
smooth + **decomposable** circuits = ...

sample variables x_1, \dots, x_n from $p(\mathbf{x})$
 \implies linear time in circuit size!



smooth + **decomposable** circuits = ...

If $p(\mathbf{x}) = \sum_i w_i p_i(\mathbf{x})$
(smoothness):

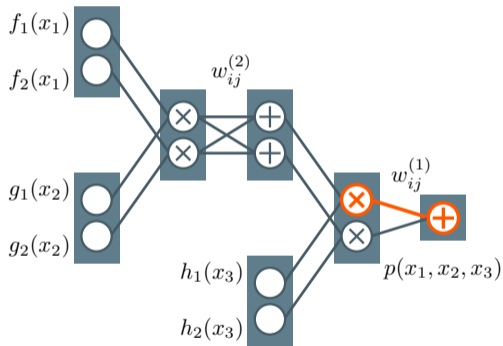


smooth + **decomposable** circuits = ...

If $p(\mathbf{x}) = \sum_i p(z = i) p_i(\mathbf{x} \mid z = i)$

(smoothness):

sample $z = i$ from $p(z)$,
then sample \mathbf{x} from $p_i(\mathbf{x} \mid z = i)$

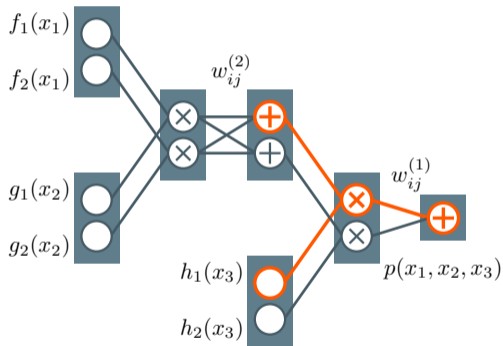


smooth + **decomposable** circuits = ...

If $p(\mathbf{x}) = p_1(\mathbf{y}) p_2(\mathbf{z})$

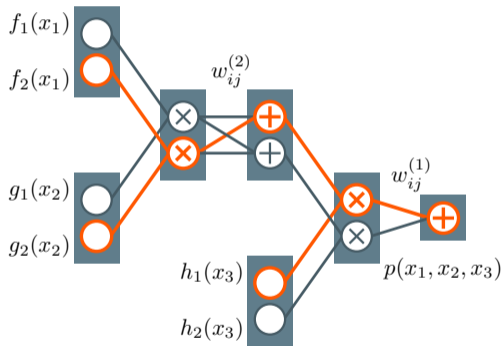
(decomposability):

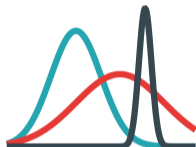
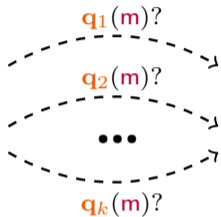
sample \mathbf{y} from p_1 and \mathbf{z} from p_2
(as they are disjoint)



smooth + **decomposable** circuits = ...

Sample from simple input distributions:
⇒ easy for Categorical, Gaussian, ...





$p_m(\mathbf{X})$

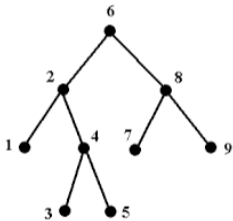
\approx

	X^1	X^2	X^3	X^4	X^5
x_8					
x_7					
x_6					
x_5					
x_4					
x_3					
x_2					
x_1					

generative models that can reason probabilistically

...but some events are certain!

When?



given \mathbf{x} // e.g. a feature map

find $\mathbf{y}^* = \operatorname{argmax}_{\mathbf{y}} p_{\theta}(\mathbf{y} \mid \mathbf{x})$ // e.g. labels of classes

s.t. $\mathbf{y} \models \mathbf{K}$ // e.g., constraints over superclasses

$$\mathbf{K} : (Y_{\text{cat}} \implies Y_{\text{animal}}) \wedge (Y_{\text{dog}} \implies Y_{\text{animal}})$$

hierarchical multi-label classification

When?



Ground Truth

given \mathbf{x} // e.g. a tile map

find $\mathbf{y}^* = \operatorname{argmax}_{\mathbf{y}} p_{\theta}(\mathbf{y} \mid \mathbf{x})$ // e.g. a configurations of edges in a grid
s.t. $\mathbf{y} \models \mathbf{K}$ // e.g., that form a valid path

// for a 12×12 grid, 2^{144} states but only 10^{10} valid ones!

nesy structured output prediction (SOP) tasks

When?



Ground Truth



ResNet-18

neural nets struggle to satisfy validity constraints!

Constraint losses



Ground Truth



ResNet-18



Semantic Loss

...but cannot guarantee consistency at test time!

SPL



Ground Truth



ResNet-18



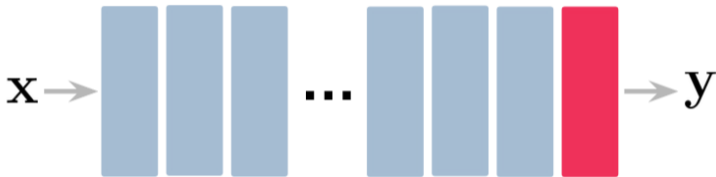
Semantic Loss



circuits

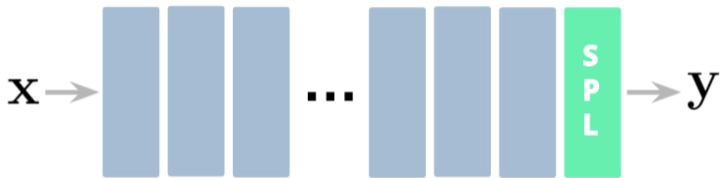
you can predict valid paths 100% of the time!

How?



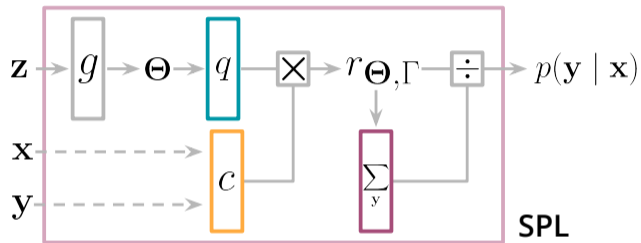
take an unreliable neural network architecture...

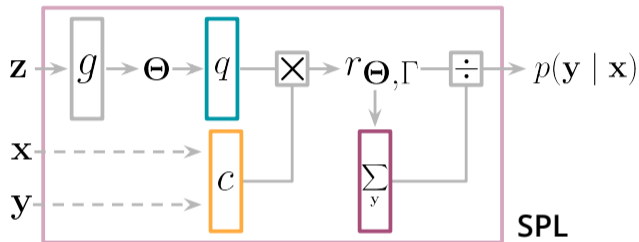
How?



*.....and replace the last layer with
a semantic probabilistic layer (SPL)*

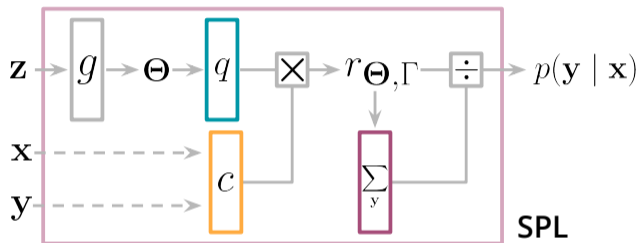
SPL





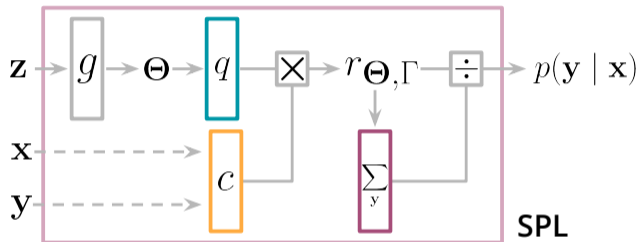
$$p(\mathbf{y} | \mathbf{x}) = \mathbf{q}_{\Theta}(\mathbf{y} | g(\mathbf{z}))$$

$\mathbf{q}_{\Theta}(\mathbf{y} | g(\mathbf{z}))$ is an expressive distribution over labels



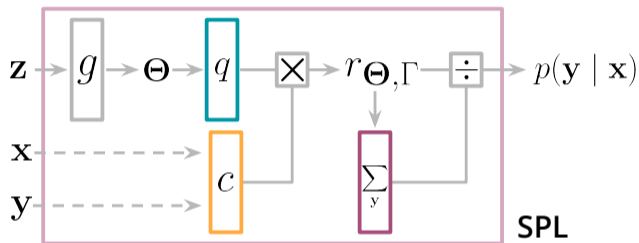
$$p(\mathbf{y} | \mathbf{x}) = \mathbf{q}_{\Theta}(\mathbf{y} | g(\mathbf{z})) \cdot \mathbf{c}_{\mathbf{K}}(\mathbf{x}, \mathbf{y})$$

$\mathbf{c}_{\mathbf{K}}(\mathbf{x}, \mathbf{y})$ encodes the constraint $\mathbb{1}\{\mathbf{x}, \mathbf{y} \models \mathbf{K}\}$



$$p(\mathbf{y} | \mathbf{x}) = \mathbf{q}_{\Theta}(\mathbf{y} | g(\mathbf{z})) \cdot \mathbf{c}_{\mathbf{K}}(\mathbf{x}, \mathbf{y})$$

a product of experts : (



$$p(\mathbf{y} | \mathbf{x}) = q_{\Theta}(\mathbf{y} | g(\mathbf{z})) \cdot c_{\mathbf{K}}(\mathbf{x}, \mathbf{y}) / \mathcal{Z}(\mathbf{x})$$

$$\mathcal{Z}(\mathbf{x}) = \sum_{\mathbf{y}} q_{\Theta}(\mathbf{y} | \mathbf{x}) \cdot c_{\mathbf{K}}(\mathbf{x}, \mathbf{y})$$

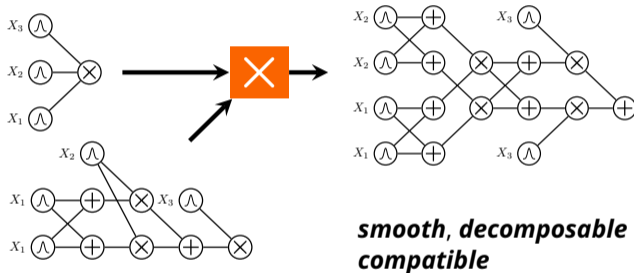
Goal

*Can we design q and c
to be **deep factorizations**
yet yielding a tractable product?*

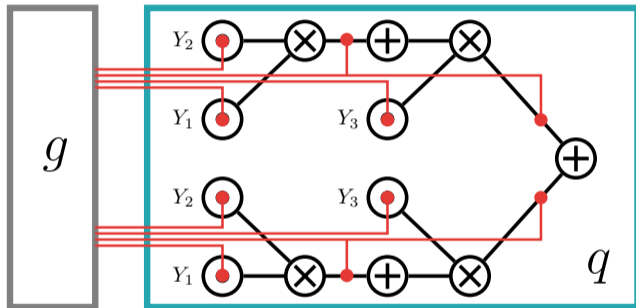
Goal

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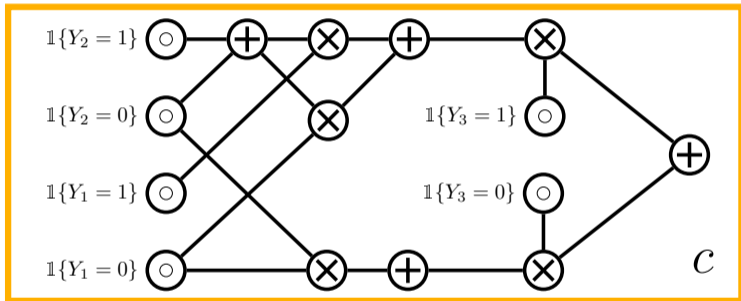
Tractable products



exactly compute \mathcal{Z} in time $O(|q||c|)$



a conditional circuit $q(y; \Theta = g(z))$



and a logical circuit $c(y, x)$ encoding K

knowledge compilation

(as a Boolean tensor factorization)

$$K: (Y_1 = 1 \implies Y_3 = 1)$$

$$\wedge (Y_2 = 1 \implies Y_3 = 1)$$

$$\text{Boolean tensor: } \mathcal{K} \in \{0, 1\}^3$$

$$k_{y_1 y_2 y_3} = \mathbf{1}\{y_1 y_2 y_3 \models K\}$$

$$\mathbf{1}\{Y_1 = 0\} \odot$$

$$\mathbf{1}\{Y_1 = 1\} \odot$$

$$\mathbf{1}\{Y_2 = 0\} \odot$$

$$\mathbf{1}\{Y_2 = 1\} \odot$$

$$\mathbf{1}\{Y_3 = 0\} \odot$$

$$\mathbf{1}\{Y_3 = 1\} \odot$$

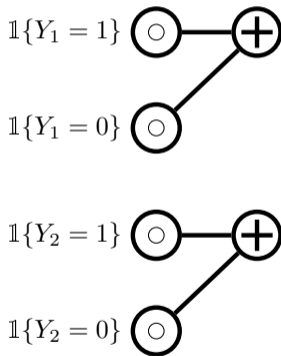
knowledge compilation

(as a Boolean tensor factorization)

$$\begin{aligned} \mathcal{K}: & (Y_1 = 1 \implies Y_3 = 1) \\ & \wedge (Y_2 = 1 \implies Y_3 = 1) \end{aligned}$$

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Pipatsrisawat and Darwiche, "New Compilation Languages Based on Structured Decomposability.", 2008

knowledge compilation

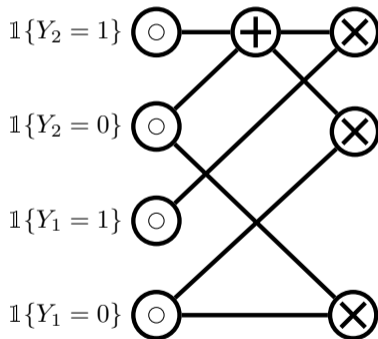
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knowledge compilation

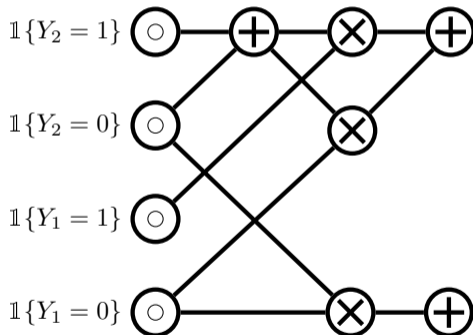
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Pipatsrisawat and Darwiche, "New Compilation Languages Based on Structured Decomposability.", 2008

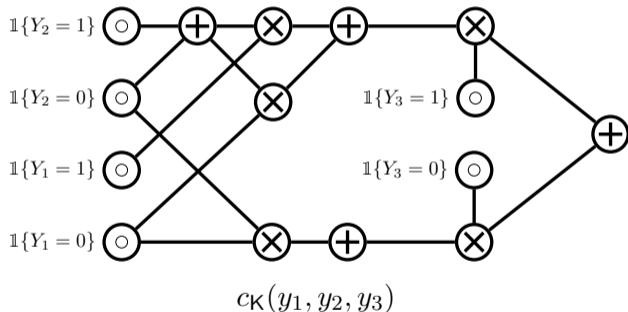
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Pipatsrisawat and Darwiche, "New Compilation Languages Based on Structured Decomposability.", 2008

Tensor Decomposition Meets Knowledge Compilation: A Study Comparing Tensor Trains with OBDDs

Ryoma Onaka, Kengo Nakamura, Masaaki Nishino, Norihito Yasuda

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more tensor factorizations for NeSy at AAI 2025

NeSy AI recipe

with circuits (and tensor factorizations)

$$K : (Y_1 = 1 \implies Y_3 = 1)$$

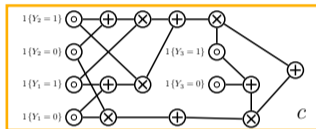
$$\wedge (Y_2 = 1 \implies Y_3 = 1)$$

1) Take a
logical constraint

NeSy AI recipe

with circuits (and tensor factorizations)

$$K : (Y_1 = 1 \implies Y_3 = 1) \\ \wedge (Y_2 = 1 \implies Y_3 = 1)$$



1) Take a
logical constraint

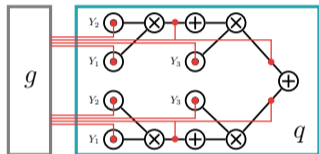
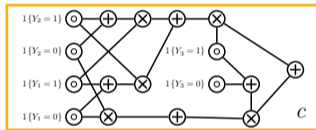
2) Compile it into
a Boolean circuit

NeSy AI recipe

with circuits (and tensor factorizations)

$$K : (Y_1 = 1 \implies Y_3 = 1)$$

$$\wedge (Y_2 = 1 \implies Y_3 = 1)$$



1) Take a logical constraint

2) Compile it into a Boolean circuit

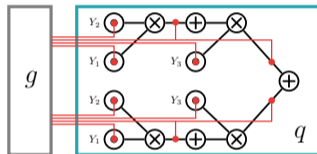
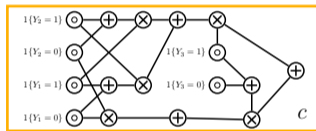
3) Multiply it by a circuit distribution

NeSy AI recipe

with circuits (and tensor factorizations)

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$$\wedge (Y_2 = 1 \implies Y_3 = 1)$$



1) Take a
logical constraint

2) Compile it into
a Boolean circuit

3) Multiply it
by a circuit distribution

4) train end-to-end by sgd!

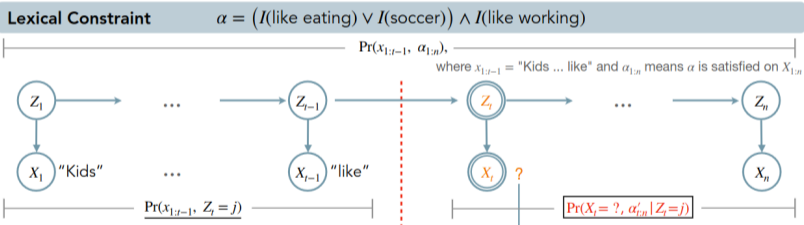
SPLS

(and more circuits)

everywhere

Tractable Control for Autoregressive Language Generation

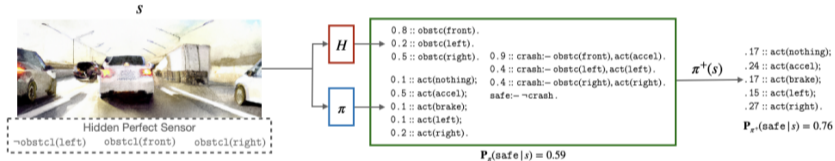
Honghua Zhang^{*1} Meihua Dang^{*1} Nanyun Peng¹ Guy Van den Broeck¹



constrained text generation with LLMs (ICML 2023)

Safe Reinforcement Learning via Probabilistic Logic Shields

Wen-Chi Yang¹, Giuseppe Marra¹, Gavin Rens and Luc De Raedt^{1,2}



reliable reinforcement learning (AAAI 23)

How to Turn Your Knowledge Graph Embeddings into Generative Models

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Nicola Di Mauro
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Robert Peharz
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enforce constraints in knowledge graph embeddings
oral at *NeurIPS 2023*

Which structural properties

for complex reasoning



q₁

$$\int p(\mathbf{x}_o, \mathbf{x}_m) d\mathbf{X}_m$$

(missing values)

smooth + decomposability



q₂

$$\mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{x}_c | X_s=0)} [f_0(\mathbf{x}_c)] - \mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{x}_c | X_s=1)} [f_1(\mathbf{x}_c)]$$

(fairness)

compatibility



q₃

$$\mathbb{E}_{\mathbf{e} \sim p_{\text{noise}}(\mathbf{E})} [f(\mathbf{x} + \mathbf{e})]$$

(adversarial robust.)

compatibility

wait...!

***“Given a reasoning task
can we **automatically** distill
a **tractable** algorithm for it?”***

A language for queries

Integral expressions that can be formed by composing these operators

$+$, \times , `pow`, `log`, `exp` and $/$

\Rightarrow *many divergences and information-theoretic queries*

A language for queries

Integral expressions that can be formed by composing these operators

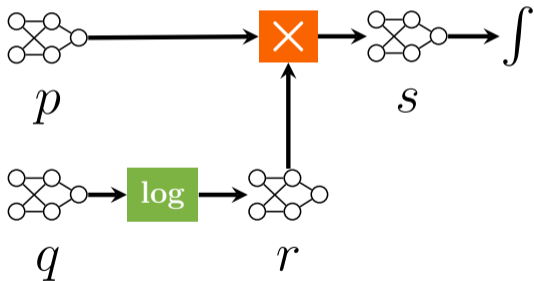
+, **×**, **pow**, **log**, **exp** and **/**

⇒ *many divergences and information-theoretic queries*

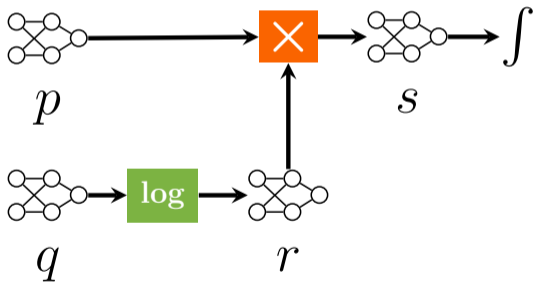
Represented as **higher-order computational graphs**—pipelines—operating over circuits!

⇒ *re-using intermediate transformations across queries*

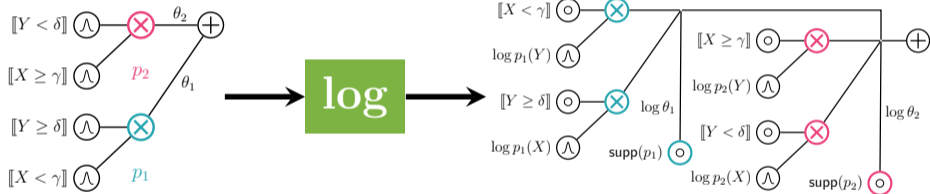
$$\text{XENT}(p \parallel q) = \int p(\mathbf{x}) \times \log q(\mathbf{x}) d\mathbf{X}$$



$$\text{XENT}(p \parallel q) = \int p(\mathbf{x}) \times \log q(\mathbf{x}) d\mathbf{X}$$



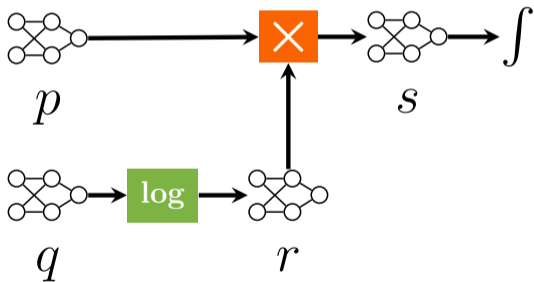
Tractable operators



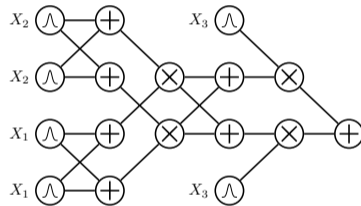
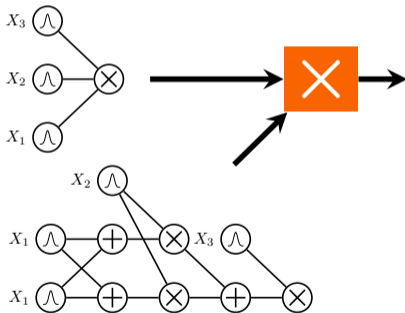
**smooth, decomposable
deterministic**

smooth, decomposable

$$\text{XENT}(p \parallel q) = \int p(\mathbf{x}) \times \log q(\mathbf{x}) d\mathbf{X}$$

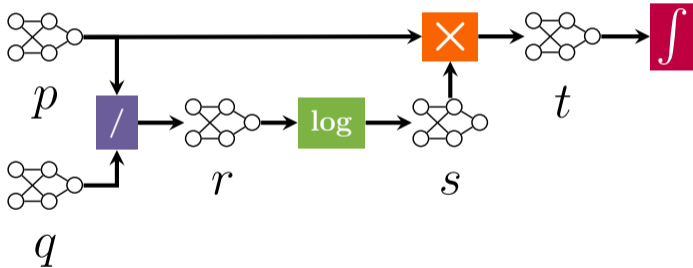


Tractable operators



***smooth, decomposable
compatible***

$$\int p(\mathbf{x}) \times \log \left(p(\mathbf{x}) / q(\mathbf{x}) \right) d\mathbf{X}$$



build a LEGO-like query calculus...

	Query	Tract. Conditions	Hardness
CROSS ENTROPY	$-\int p(\mathbf{x}) \log q(\mathbf{x}) d\mathbf{X}$	Cmp, q Det	#P-hard w/o Det
SHANNON ENTROPY	$-\sum p(\mathbf{x}) \log p(\mathbf{x})$	Sm, Dec, Det	coNP-hard w/o Det
RÉNYI ENTROPY	$(1 - \alpha)^{-1} \log \int p^\alpha(\mathbf{x}) d\mathbf{X}, \alpha \in \mathbb{N}$	SD	#P-hard w/o SD
MUTUAL INFORMATION	$(1 - \alpha)^{-1} \log \int p^\alpha(\mathbf{x}) d\mathbf{X}, \alpha \in \mathbb{R}_+$	Sm, Dec, Det	#P-hard w/o Det
KULLBACK-LEIBLER DIV.	$\int p(\mathbf{x}, \mathbf{y}) \log(p(\mathbf{x}, \mathbf{y}) / (p(\mathbf{x})p(\mathbf{y})))$	Sm, SD, Det*	coNP-hard w/o SD
RÉNYI'S ALPHA DIV.	$\int p(\mathbf{x}) \log(p(\mathbf{x}) / q(\mathbf{x})) d\mathbf{X}$	Cmp, Det	#P-hard w/o Det
ITAKURA-SAITO DIV.	$(1 - \alpha)^{-1} \log \int p^\alpha(\mathbf{x}) q^{1-\alpha}(\mathbf{x}) d\mathbf{X}, \alpha \in \mathbb{N}$	Cmp, q Det	#P-hard w/o Det
CAUCHY-SCHWARZ DIV.	$(1 - \alpha)^{-1} \log \int p^\alpha(\mathbf{x}) q^{1-\alpha}(\mathbf{x}) d\mathbf{X}, \alpha \in \mathbb{R}$	Cmp, Det	#P-hard w/o Det
SQUARED LOSS	$\int [p(\mathbf{x}) / q(\mathbf{x}) - \log(p(\mathbf{x}) / q(\mathbf{x})) - 1] d\mathbf{X}$	Cmp, Det	#P-hard w/o Det
	$-\log \frac{\int p(\mathbf{x}) q(\mathbf{x}) d\mathbf{X}}{\sqrt{\int p^2(\mathbf{x}) d\mathbf{X} \int q^2(\mathbf{x}) d\mathbf{X}}}$	Cmp	#P-hard w/o Cmp
	$\int (p(\mathbf{x}) - q(\mathbf{x}))^2 d\mathbf{X}$	Cmp	#P-hard w/o Cmp

...and **compositionally** derive many more tractable algorithms

Takeaways

1 *faster sampling* routines for non-negative factorizations

Takeaways

- 1 *faster sampling* routines for non-negative factorizations
- 2 integration of *logical constraints* with guarantees

Takeaways

- 1 *faster sampling* routines for non-negative factorizations
- 2 integration of *logical constraints* with guarantees
- 3 *automatically distill* efficient algorithms for tensor networks via circuit properties

Takeaways

Questions?

- 1 *faster sampling* routines for non-negative factorizations
- 2 integration of *logical constraints* with guarantees
- 3 *automatically distill* efficient algorithms for tensor networks via circuit properties

outline

- 1 connecting *tensor factorizations* and *circuits*
- 2 a *unifying pipeline* to build factorizations & circuits
- 3 a *property-driven* approach to inference & reasoning
- 4 *expressiveness analysis*: known and new results

tl;dr

***“Understand when and how
one factorization scheme can be
provably more expressive than others”***

Expressiveness of tensor factorizations

We care about factorization methods that yield **compact** decompositions
(minimise memory footprint & computation)

Expressiveness of tensor factorizations

We care about factorization methods that yield **compact** decompositions
(minimise memory footprint & computation)

*“if rank(s) is **exponential in** d , then it is not useful!”*

\implies storing $\mathcal{T} \in \mathbb{R}^{M \times \dots \times M}$ requires $\mathcal{O}(M^d)$ memory

Expressiveness of tensor factorizations

We care about factorization methods that yield **compact** decompositions
(minimise memory footprint & computation)

*“if rank(s) is **exponential in** d , then it is not useful!”*

\implies storing $\mathcal{T} \in \mathbb{R}^{M \times \dots \times M}$ requires $\mathcal{O}(M^d)$ memory

One factorization method may require **exponentially smaller rank** than others \implies it is more **expressive**

wait...!

what about circuits?

Expressiveness of circuits

A rigorous concept in circuit complexity theory

Valiant, "Negation can be exponentially powerful", 1979

Darwiche and Marquis, "A knowledge compilation map", 2002

Martens and Medabalimi, "On the expressive efficiency of sum product networks", 2014

Expressiveness of circuits

A rigorous concept in circuit complexity theory

Expressiveness results of circuits based on the **circuit size**

⇒ number of edges between units (amount of computation)

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Expressiveness of circuits

A rigorous concept in circuit complexity theory

Expressiveness results of circuits based on the **circuit size**

⇒ number of edges between units (amount of computation)

Different circuit classes have different expressive power

Valiant, "Negation can be exponentially powerful", 1979

Darwiche and Marquis, "A knowledge compilation map", 2002

Martens and Medabalimi, "On the expressive efficiency of sum product networks", 2014

“Circuit complexity theory helps proving stronger results for tensor factorizations”

SUBTRACTIVE MIXTURE MODELS VIA SQUARING: REPRESENTATION AND LEARNING

Lorenzo Loconte^{1*} **Aleksanteri M. Sladek²** **Stefan Mengel³**
Martin Trapp² **Arno Solin²** **Nicolas Gillis⁴** **Antonio Vergari¹**

¹ School of Informatics, University of Edinburgh, UK

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³ University of Artois, CNRS, Centre de Recherche en Informatique de Lens (CRIL), France

⁴ Department of Mathematics and Operational Research, Université de Mons, Belgium

Accepted at ICLR 2024 as a spotlight

Monotonic probabilistic circuits

Monotonic circuits

$$p(\mathbf{x}) = \frac{1}{Z} c(\mathbf{x}), \quad c(\mathbf{x}) \geq 0$$

where parameters and input functions are **positive**

Monotonic probabilistic circuits

Monotonic circuits

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where parameters and input functions are **positive**

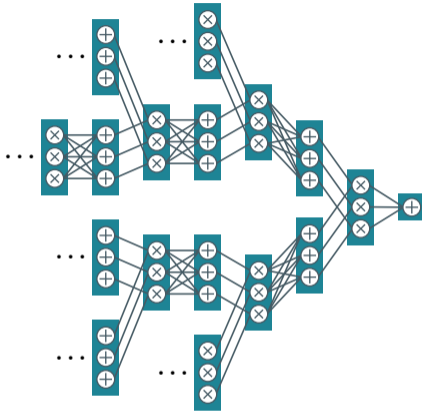
(represent non-negative tensor factorizations *[Cichocki and Phan 2009]*)

A limitation of monotonic circuits

 = set of distributions modeled by **polysize** circuits



- UDISJ



$\exists p$ requiring **exponentially large monotonic circuits...**

Squared circuits

$$p(\mathbf{x}) = \frac{1}{Z} c^2(\mathbf{x}), \quad c(\mathbf{x}) \in \mathbb{R}$$

where parameters and input functions can be **negative**

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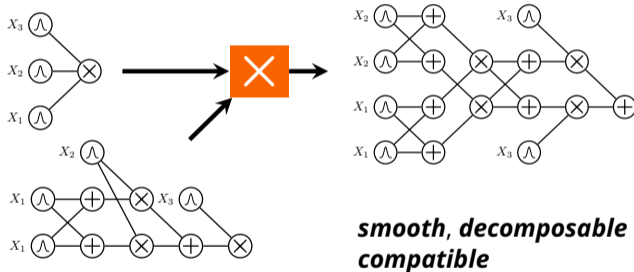
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Tractable product

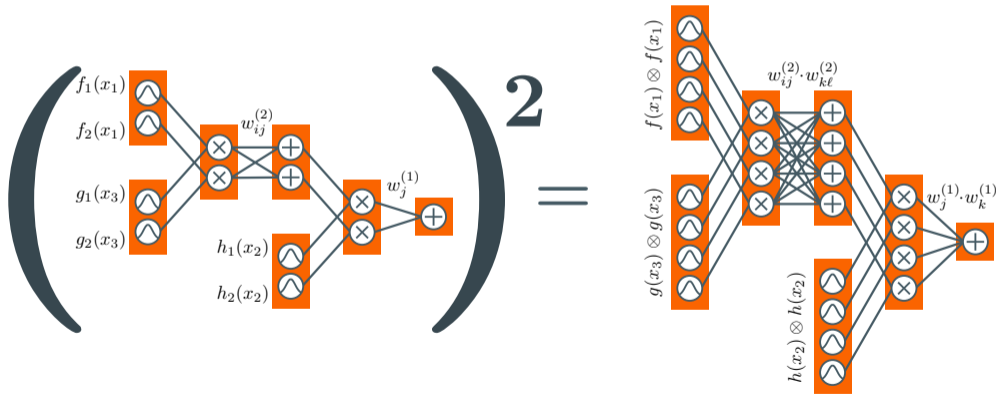
thanks to circuit compatibility

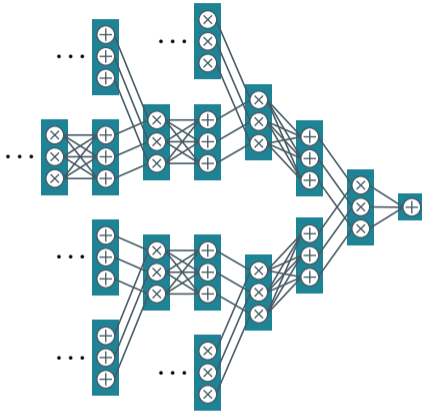


Squared circuits

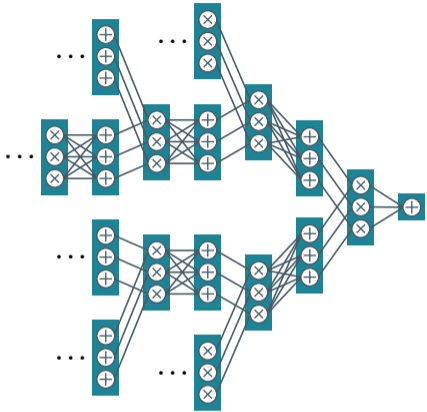
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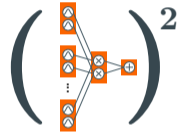
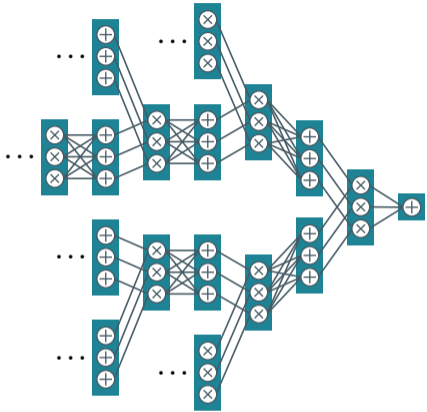


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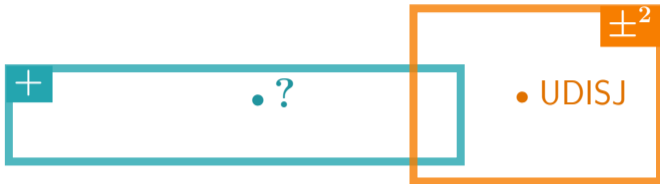


$$\left(\begin{array}{c} \text{+} \\ \text{+} \\ \vdots \\ \text{+} \end{array} \right)^2$$

...instead **squared circuits** require **polynomial size**

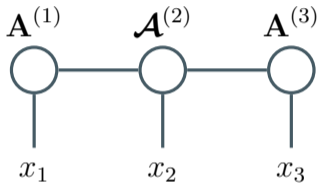


Squared circuits more expressive than **monotonic** ones



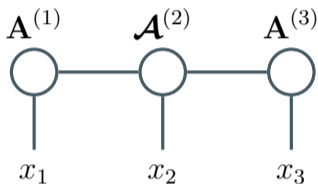
“Can monotonic circuits be more expressive than squared?”

Squared circuits are sparse Born machines

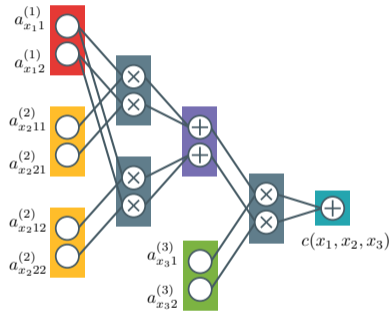


$$t_{x_1 x_2 x_3} = \sum_{r_1=1}^R \sum_{r_2=1}^R a_{x_1 r_1}^{(1)} a_{r_1 x_2 r_2}^{(2)} a_{r_2 x_3}^{(3)}$$

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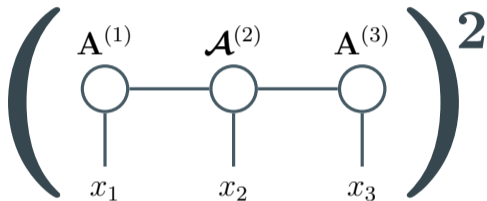


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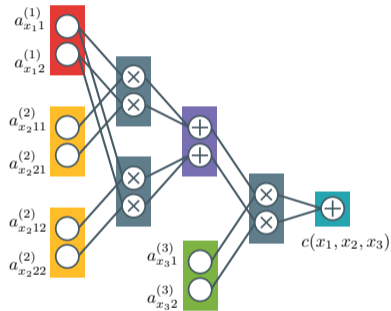
circuit compatible with itself

Squared circuits are sparse Born machines



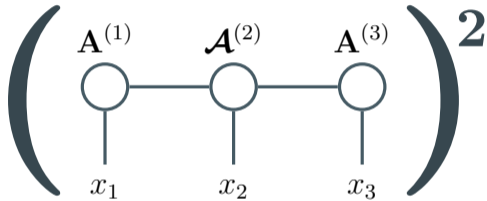
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$$p(x_1, x_2, x_3) \propto (t_{x_1 x_2 x_3})^2 \quad (\text{Born machine})$$



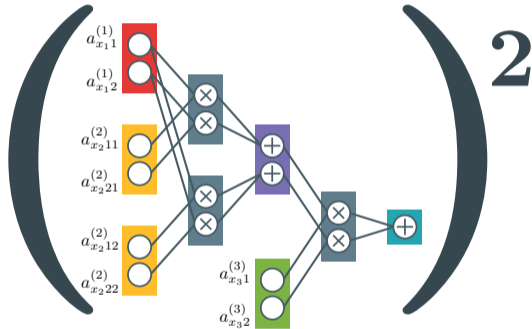
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A limitation of Born machines

(with real tensor-train factorization)

Proposition 5 (Glasser et al. 2019)

There exists non-negative tensors over $2d$ variables that can be factorized as **positive TT of constant rank 2**, but **real Born machines** have at least rank $2^{\Omega(d)}$.

A limitation of Born machines

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Proposition 5 (Glasser et al. 2019)

There exists non-negative tensors over $2d$ variables that can be factorized as **positive TT of constant rank 2**, but **real Born machines** have at least rank $2^{\Omega(d)}$.

“Can it be generalized to squared circuits?”



“Can monotonic circuits be more expressive than squared?”

Yes!

Sum of Squares Circuits*

Lorenzo Loconte  Stefan Mengel  Antonio Vergari 

 School of Informatics, University of Edinburgh, UK

 Univ. Artois, CNRS, Centre de Recherche en Informatique de Lens (CRIL), France

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Accepted at AAAI 2025

Poster #840, Hall E, Thursday February 27,
12:30pm-2:30pm

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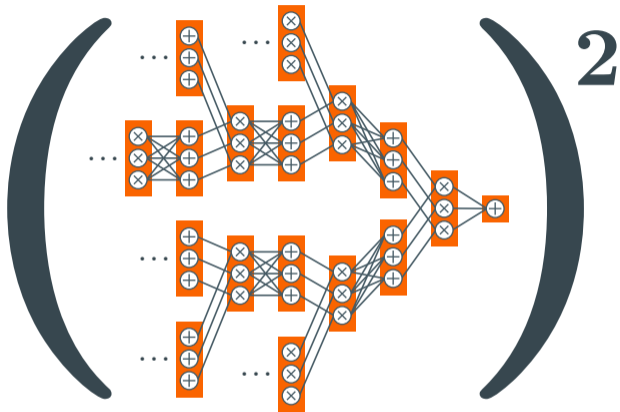
Wang and Van den Broeck

“On the Relationship Between Monotone
and Squared Probabilistic Circuits”

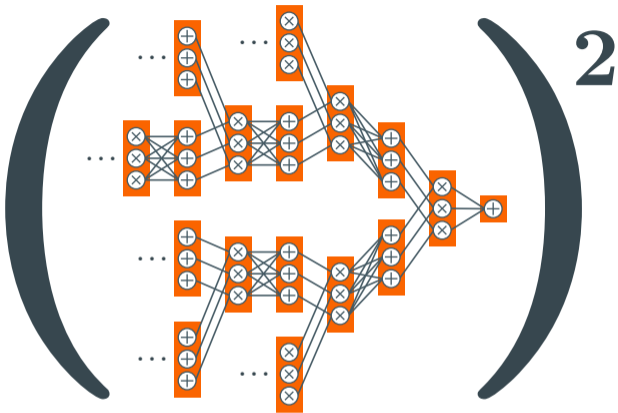
(also at AAI 2025)



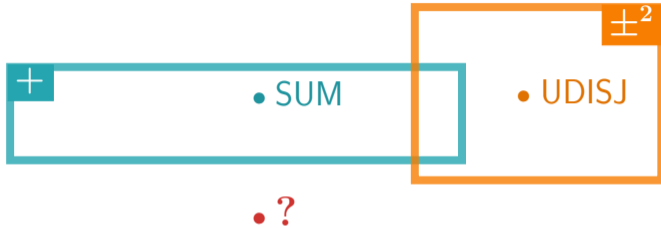
$\exists p$ requiring **polysize monotonic circuits...**



...but require **exponentially large squared circuits**



Squaring alone can reduce expressiveness!
 (generalizes to factorizations other than tensor-trains)



“How to build circuits more expressive than both?”

Sum of squares (SOS) circuits

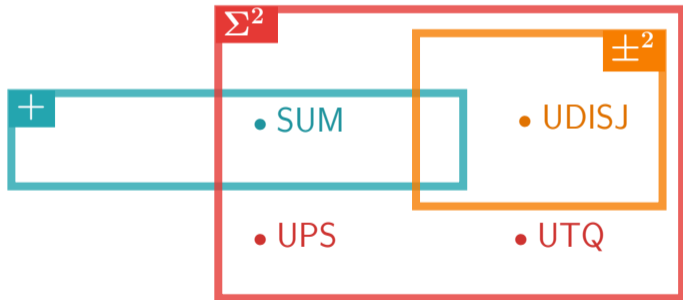
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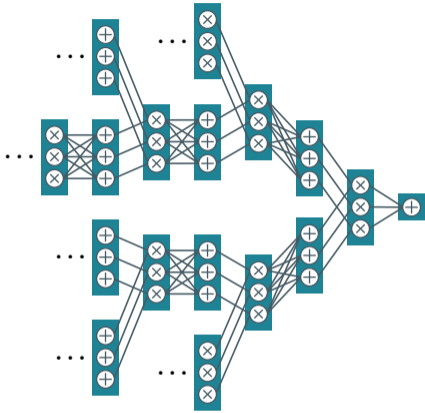
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Sum of squares (SOS) circuits

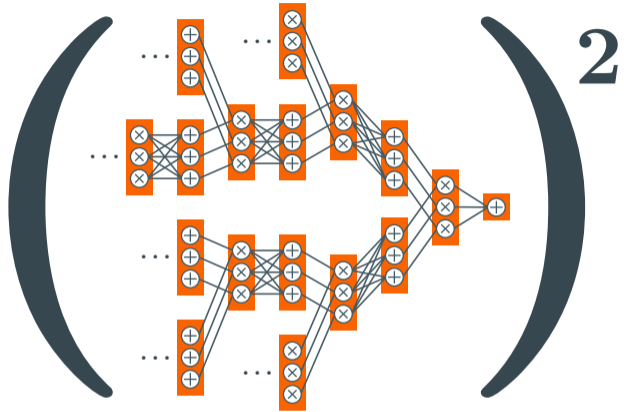
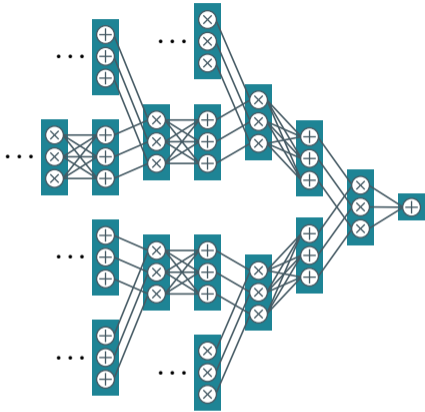
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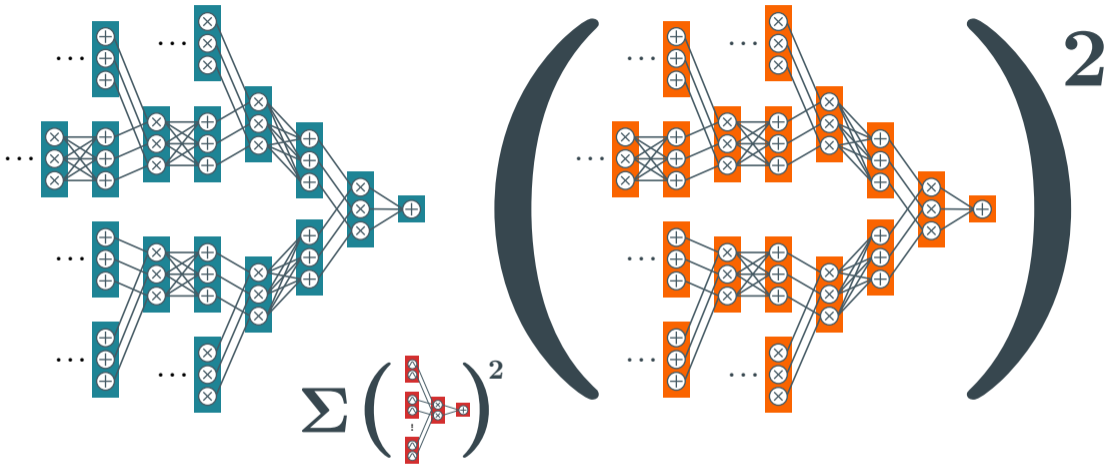




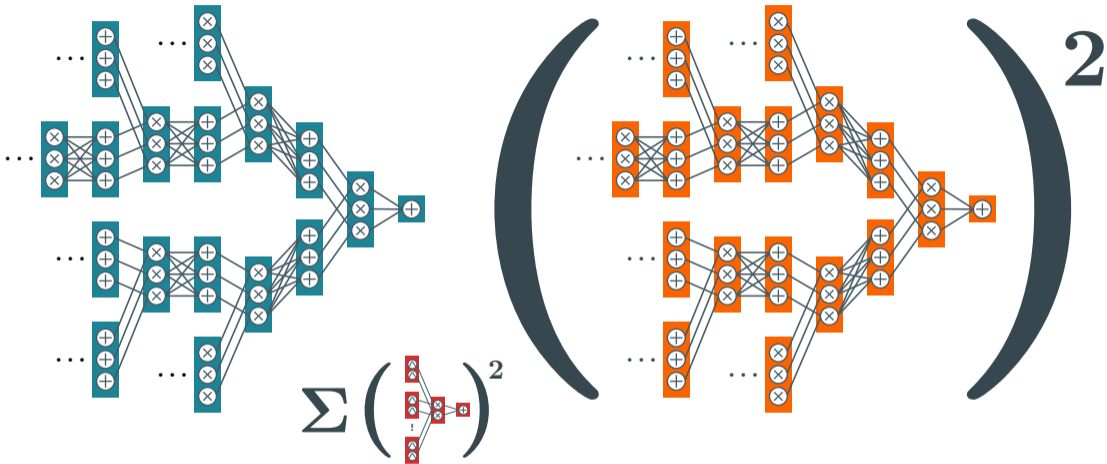
$\exists p$ requiring **exponentially large monotonic circuits...**



...and also **exponentially large squared circuits** ...

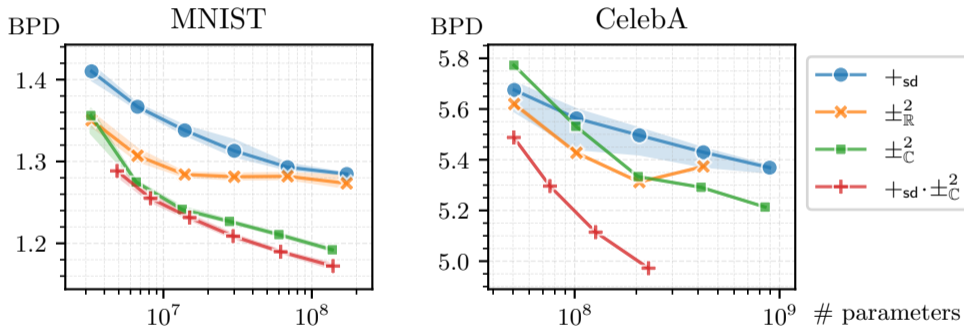


...but a **sum of squares (SOS)** polysize circuits



SOS can surpass both expressiveness limitations!

Experiments



Complex squared circuits are SOS (and scale better!)

Takeaways

- 1 factorizations and circuits **expressiveness results...**
bridge rank and circuit size

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Takeaways

Questions?

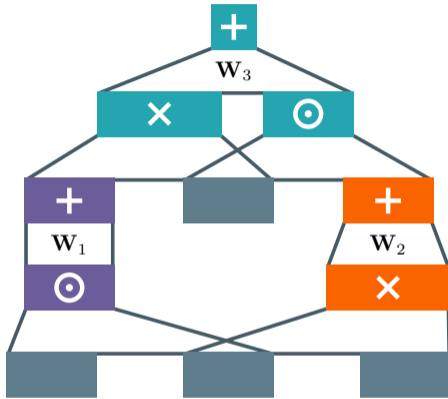
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wait...!

conclusions...?

PC ARCHITECTURE	REGION GRAPH	SUM-PRODUCT LAYER	FOLD
Poon&Domingos (Poon & Domingos, 2011)	PD	CP^T	✗
RAT-SPN (Peharz et al., 2020c)	RND	Tucker	✗
EiNet (Peharz et al., 2020a)	{ RND, PD }	Tucker	✓
HCLT (Liu & Van den Broeck, 2021b)	CL	CP^T	✓
HMM/MPS $_{\mathbb{R} \geq 0}$ (Glasser et al., 2019)	LT	CP^T	✗
BM (Han et al., 2018)	LT	CP^T	✗
TTDE (Novikov et al., 2021)	LT	CP^T	✗
NPC ² (Loconte et al., 2024)	{ LT, RND }	{ CP^T , Tucker }	✓
TTN (Cheng et al., 2019)	QT-2	Tucker	✗
Mix & Match (our pipeline)	{ RND, PD, LT, CL, QG, QT-2, QT-4 }	{ Tucker, CP, CP^T } \cup { CP^S , CP^{XS} FOLD ✓ }	{ ✗, ✓ }

takeaway #1: unifying a fragmented literature



takeaway #2: easily build novel factorizations



Ground Truth



ResNet-18

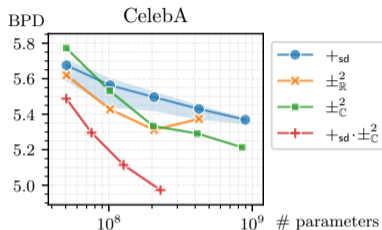
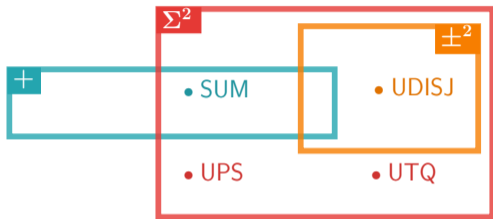


Semantic Loss



circuits

takeaway #3: use them for efficient & reliable inference



takeaway #4: SOS circuits are provably more expressive factorizations

What is the Relationship between Tensor Factorizations and Circuits (and How Can We Exploit it)?

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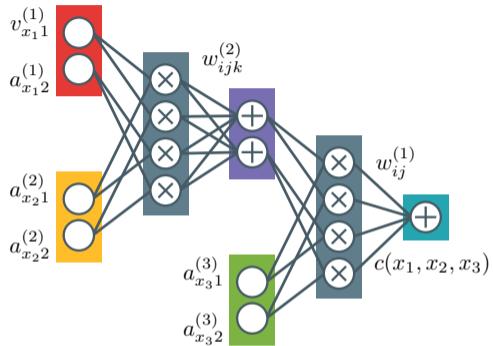
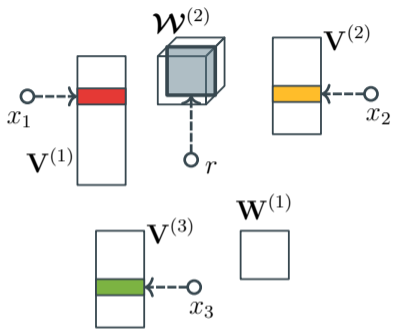
learning & reasoning with circuits in pytorch

`https://github.com/april-tools/circuit`

colorai
connecting
low-rank
representations
in ai

workshop at AAAI-25, March 4

`april-tools.github.io/colorai/`



questions?